

METHOD AND APPARATUS FOR COMPOSITE USER INTERFACE CREATION

Field of Invention

The present invention relates to methods and systems for generating composite user interfaces, and to methods and systems for directing messages between a composite user interface and a plurality of source applications.

Background of Invention

Computer users routinely need to use a plurality of different applications in order to complete tasks allocated to them, and each application typically has a separate user interface. Switching between the different user interfaces of the different applications in order to complete a given task considerably degrades user efficiency. It will often be the case that different applications are supplied by different vendors and accordingly their user interfaces have a different “look and feel”, further degrading operator efficiency.

For example, in order to process customer enquiries, operators in a call centre may need to access a customer management application to access customer details, a billing application to access customer account information, and a payment application to process any payment which may be made by the customer over the telephone, for example by credit card. Working in this manner is inefficient, given that the operator is required to switch between applications in order to complete some tasks. Furthermore, a customer will typically remain on the telephone while the operator uses these different applications, and it is therefore advantageous to speed up the processing of enquires, in order to offer a higher quality customer service.

Various proposals have been made to enhance user efficiency when multiple applications need to be used.

The multiple applications can be combined into a single product or product suite. While such a proposal provides great increases in user efficiency, it is difficult and expensive

to implement. Furthermore, such a combined product or product suite will typically have a different user interface from those used previously, therefore meaning that users need to be trained in use of the combined product, further increasing cost.

It has alternatively been proposed that the multiple application can be combined in some way. For example, all requests can be passed to a single one of the applications, and this application can be adapted to forward requests to an appropriate source application. Such a solution typically requires considerable customisation if it is to work in under all circumstances that may routinely arise, making such a solution difficult to implement.

It is an object of the present invention to obviate or mitigate at least some of the problems outlined above.

Summary of Invention

In accordance with one aspect, the present invention provides a method and system for processing messages within a computer system. A message comprising details of a predetermined set of services and a command is received, and the command is executed. The message is then transmitted to a service in said predetermined set of services. In accordance with some embodiments of the invention, a plurality of services is provided and each service executes a respective command specified by the message. Each service in the predetermined set of services uses said details to determine whether the message should be sent to another service. If it is determined that the message should be sent to another service the message is then transmitted to an appropriate service.

The present invention thus provides a distributed messaging method, in which message routing is determined by data contained within individual messages, and is achieved by the services to which they are directed. The invention therefore removes the need for any central messaging service which is responsible for all message routing operations. In order to allow distributed routing, the message may comprise an ordered list of pairs, a first element of each pair representing a service in the predetermined set of services, and a second element of each pair representing a command to be executed by that service, thus each service can determine a command to be executed and can determine a next service to which the message should be directed.

The term service is to be understood broadly to cover any processing means or processing module which is adapted to receive a message, carry out processing specified by that message, and forward the message to a further service using data specified within the message. The services can be implemented using computer program code means. For example each service can be implemented as an instance of a class defined in an object oriented programming language such as Java or C++. Each class representing a service preferably has an associated service handler class which specifies a method configured to execute a command directed to an instance of the respective service class. Advantageously, each service object references a plurality of service handler objects which are instances of the respective service handler class. By providing a plurality of service handler objects in this way, a single service may carry out a plurality of commands concurrently, thereby improving processing efficiency.

In one embodiment, a plurality of services in the predetermined set of services may operate within a single operating system process, whilst in other embodiments of the invention some services in the predetermined set of services may operate within a plurality of operating system processes. Providing a plurality of processes enhances scalability, while operating a plurality of services within a single operating system process can result in improved performance. It is preferred that an attempt is made to locate a service within the current operating system process, and if such an attempt is successful, relatively costly inter process communication can be avoided. However, if the attempt is unsuccessful, the message may be transmitted to a messaging service within the current operating system process which is responsible for inter-process communication, and the messaging service may transmit the message to a different operating system process, thus providing scalability. For example, in some embodiments, the system and method of the invention allow different services provided on different computers connected to a computer network to be handled in a similar manner to different services provided within different operating system processes on a single computer. It is preferred that messaging is achieved using the Java Messaging Service (JMS).

In some embodiments, a message may be directed between a composite user interface and at least one source application, such that the composite user interface can be used to interface with the at least one source application. A message emanating from a composite user interface may be directed to a service which generates at least one further message. The further message may comprise details of a further set of services to which the further message is to be directed, and the further message can then be directed in the distributed manner described above. The at least one further message may be processed by one of services in said further set of services to produce a request which is transmitted to the at least one source application. Messaging methods provided by embodiments of the present invention offer particular benefits in composite user interface applications given that communication between a composite application and one or more source applications can be provided in a decoupled manner. Thus, the invention allows composite user interfaces to be conveniently and efficiently provided allowing users to use a composite user interface to control a plurality of source applications, offering considerable benefits in user productivity.

Data from the source application may be used to create a response message in response to said request. The response message can again be handled in the distributed manner described above. An aggregation service in the set of response services may receive a plurality of response messages and may combine said plurality of response messages to create a further response message. Thus, the aggregation service can effectively combine a plurality of messages containing user interface elements to create a composite user interface.

In some embodiments, the aggregation service may also generate additional user interface elements which are combined with said user interface to generate said composite user interface. Adding additional user interface elements in this way can be useful in providing a unified “look and feel” throughout a composite application.

The aggregation service may combine said user interface elements to generate the composite user interface in accordance with predefined configuration data, which can suitably be stored in a hierarchical data structure. A first entity within the hierarchical data structure may represent the composite user interface, and child entities of said first

entity may represent the plurality of source user interface elements for inclusion in the composite interface. Each child entity of the first entity can in turn have child entities and so on. Using a hierarchical data structure is advantageous as it allows easy scalability where user interface elements are nested within other user interface elements.

In some embodiments, each child entity has an associated parameter indicating whether the respective source user interface element is mandatory. The aggregation service may receive a user interface element, store data indicative of receipt of said user interface element in a data structure associated with said hierarchical data structure, and when all source user interface elements having an associated parameter indicating that the source user interface element is mandatory have been received, combine said plurality of source data items to generate said composite user interface. Differentiating between mandatory and non-mandatory user interface elements in this way can provide considerable benefits in terms of efficiency, and allows a user interface to be provided as soon as key components have been received.

The composite user interface may be generated in an internal format, which may be converted to an output format using output configuration data. This allows the invention to be easily adapted for use with a plurality of different output formats, assuming only that appropriate configuration data is provided.

A transformation service may receive data from the at least one source application and transform said data into an internal format, said transformed data being contained in said plurality of response messages. Said transformation may be effected in a plurality of different ways, for example using regular expressions, or using a class defined in an object oriented programming language to transform data.

In some embodiments, the aggregation service may be configured to expect to receive a predetermined number of response messages in response to transmission of said further message, and said further response message may be generated when the predetermined number of response messages has been received.

According to a further aspect of the present invention, there is provided a method and system for generating a composite user interface for presentation to a user. The method comprises generating requests for a plurality data items for inclusion in the interface, transmitting each request to one of a plurality of source applications, and combining data items received in response to at least one of said requests to generate the user interface. At least some of the predetermined plurality of data items are mandatory, and at least some of the predetermined plurality of data items are optional, and the composite user interface is generated when all mandatory data items have been received. This aspect of the present invention allows composite user interfaces to be generated more quickly, given that it is necessary to await receipt only of mandatory user interface elements.

The method preferably comprises generating said plurality of requests from a single request such as a HTTP request entered by a user using a web browser. The plurality of requests are decoupled from the single request. Said single request may be received by an aggregation service functioning in a manner similar to that described above.

A further aspect of the present invention, provides a method for generating a composite user interface for presentation to a user, said composite user interface comprising a plurality of user interface elements generated from source interface elements provided by at least one source application. The method comprises generating a plurality of request messages, transmitting each request message to an appropriate source application and receiving a plurality of source interface elements from the at least one source application. Source interface elements are compared with a plurality of predefined source interface templates; and if said received source interface element matches a predefined source interface template, at least one user interface element is extracted for inclusion in said composite user interface. At least one of said source interface elements may be a HTML document, or a part of a HTML document such as, for example, a text box, fixed text data, or buttons responsive to selection by a pointing device.

The method may further comprise: creating an internal representation of each extracted user interface element, forwarding said internal representations to an aggregation

service, and combining said internal representations to create an internal representation of the composite user interface. The method may further comprise transforming said internal representation of the composite user interface into an output format specified by predefined configuration data. By using internal representations in this way, the process of composite user interface creation can be carried out independently of any required output formats, meaning that the method can be easily applied to a wide variety of different output formats (e.g. HTML for a web based output), assuming only that necessary configuration data is supplied.

The invention also provides a carrier medium such as a disk or CD-ROM carrying computer readable program code means for controlling a computer to carry out any of the methods described above.

The invention also provides a computer apparatus comprising a program memory containing processor readable instructions, and a processor for reading and executing the instructions contained in the program memory. The processor readable instructions comprising instructions controlling the processor to carry out any of the methods described above.

Brief description of drawings

Embodiments of the present invention will now be described, by way of example, with reference to the accompanying drawings, in which:

Figure 1 is a schematic overview of a system for composite application creation in accordance with an embodiment of the present invention;

Figure 2 is a schematic illustration showing the system of Figure 1 in further detail;

Figure 3 is a schematic illustration of process management within the server of Figure 2;

Figure 4 is a schematic overview of Java classes used to implement some parts of the architecture of Figure 3;

Figures 4A to 4V are UML class diagrams showing the classes of Figure 4 in further detail;

Figure 5 is a schematic overview of Java classes used to implement the nodes illustrated in Figure 3;

Figures 5A to 5J are UML class diagrams showing the classes of Figure 5 in further detail;

Figure 6 is a state transition diagram showing transitions between classes of Figure 5 used to represent state information related to nodes;

Figure 7 is a schematic overview of Java classes used to implement the services illustrated in Figure 3;

Figures 7A to 7F are UML class diagrams showing the classes of Figure 7 in further detail;

Figure 8 is a schematic overview of Java classes class used to implement messaging within the architecture of Figure 3;

Figures 8A to 8K are UML class diagrams showing the classes of Figure 8 in further detail;

Figure 9 is a schematic illustration showing creation of a process within the architecture of Figure 3;

Figure 10 is a schematic illustration showing creation of a node within one of the processes of Figure 3;

Figure 11A is a schematic overview of a messaging process in accordance with an embodiment of the present invention;

Figure 11B is a flowchart showing the messaging process of Figure 11A in further detail;

Figures 12A to 12D are schematic illustrations showing the messaging of Figures 11A and 11B in further detail;

Figure 13 is a schematic illustration showing distributed messaging as illustrated in Figures 12A to 12D;

Figure 14 is a schematic illustration of the logical architecture of the web server and the server of Figure 2;

Figure 15 is a schematic illustration showing how parts of the architecture of Figure 14 cooperate to provide composite user interfaces;

Figure 16 is a table showing invocation parameters used by the web server of Figure 2;

Figure 17 is a table showing invocation parameters used by the server of Figure 2;

Figure 18 is a table showing Hypertext Markup Language (HTML) tags which may be used in embodiments of the present invention;

Figure 19 is an extract from a configuration file, illustrating how various parameters can be initialised;

Figure 20 is a tree diagram illustrating configuration data for the server and web server of Figure 2;

Figure 21 is a tree diagram showing configuration data pertinent to the Communications Layer (CL) services of Figure 14;

Figures 22 and 23 are tree diagrams showing configuration data pertinent to the Data Transformation (DT) services of Figure 14;

Figures 24, 25 and 26 are tree diagrams showing configuration data pertinent to the User Experience Manager (UXM) service of Figure 14;

Figure 27 is a tree diagram showing configuration data pertinent to the Application Lockdown Framework (ALF) illustrated in Figure 14;

Figure 28 is a tree diagram showing configuration data for the web server illustrated in Figure 14;

Figure 29 shows configuration data for the IDM illustrated in Figure 14;

Figure 30 is a schematic illustration showing a relationship between a UXM tree and a UXMObject tree;

Figure 31 is a schematic illustration showing how a HTML document may be converted into a plurality of UXM objects;

Figure 32 is a schematic illustration of operation of the DT service for a particular source application; and

Figure 33 is a HTML code fragment which can be used to generate an IncomingUserRequest message.

Description of preferred embodiments

Figure 1 illustrates a system for creating composite applications in accordance with an embodiment of the present invention. A first source application provides a first user interface 1, a second source application provides a second user interface 2, and a third source application provides a third user interface 3. A composition system 4 composes the first, second and third user interfaces to form a first composite user interface 5 and a

second composite user interface 6. It can be seen that the first composite user interface 5 is created from the composition of the first, second and third user interfaces 1, 2, 3, while the second composite user interface 6 is created from the composition of the second and third user interfaces 2, 3.

The composition system 4 processes requests from users of the composite user interfaces 5, 6 and generates requests to the appropriate source applications. The composition system 4 additionally receives information from the first, second and third applications in response to such requests, and uses this information to generate the composite user interfaces 5, 6.

For example, the composite user interfaces 5 may comprise a number of pages each containing elements from one or more of the first, second and third user interfaces 1, 2, 3. That is a first page of the user interface may be made up solely from the first user interface 1, and a second page may be up from the combination of part of the second user interface 2 and part of the third user interface 3, and a third page may comprise part of the second user interface 2, rearranged by the composer to provided a “look and feel” which is similar to that of the first and second pages.

Figure 2 illustrates the system of Figure 1 in further detail. The composite user interface 5 is displayed to a user by means of a display device 7 connected to a PC 8. Similarly, the composite user interface 6 is displayed to a user by means of a display device 9 connected to a PC 10. The PCs 8, 10 have means for connection to the Internet 11. The PCs 8, 10 can be connected to the Internet 11 via a local area network connection, or alternatively via a modem (not shown).

A Web server 12 is also connected to the Internet 11, and stores a plurality of webpages in HTML (Hypertext Markup Language) format which can be accessed by the PCs 8, 10. The Web server is connected to a server 13. This connection can either be a direct connection 14 (as shown in Figure 2), or alternatively a connection across a network such as the Internet. The server 13 is in turn connected to three servers 15, 16, 17 which respectively execute the first second and third applications to provide the first second

and third user interfaces 1, 2 3. The server 13 is also connected to a database 18 which stores configuration data.

In operation, the web server 12 provides HTML documents which form the composite user interfaces 5, 6. The composite user interfaces 5, 6 are created by processes running on the server 13, which create the composite user interfaces from information provided by the servers 15, 16, 17, in accordance with configuration data stored in the database 18 as described below. Data input by users of the composite interfaces 5, 6 is received by the web server 12 and forwarded to the server 13 via the connection 14. The server 13 processes such data in a predefined manner, and forwards data to the servers 15, 16, 17 as appropriate.

It will be appreciated that in some embodiments of the present invention, instead of connection over the Internet 11, the appropriate connections can be realised through an organisation's Intranet.

Operation of the Web server 12 and the server 13 in creating and modifying the composite user interfaces 5, 6 is described below.

The server 13 operates using a flexible process framework. Figure 3 illustrates components used within this flexible process framework.

A HostMasterWatchdog component 19 checks if a HostMaster process 20 is already running, and if not causes a HostMaster process to be started. The HostMaster process 20 is responsible for creating and removing processes. The HostMaster process 20 ensures that all created processes are up and running and also passes relevant configuration information to such processes. Any host operating in accordance with the process framework described herein runs a single instance of the HostMasterWatchdog component 19 and the HostMaster process 20.

In the illustration of Figure 3, a server runs a first process 21 and a second process 22. The first process comprises two nodes 23, 24 and the second process comprises two nodes 25, 26. Nodes are containers for user services, and are described in further detail

below. Each of the processes 21, 22 includes a respective process agent 27, 28 which is responsible for starting nodes within the respective process, and restarting nodes within the process in case of failure. Operation of nodes within the processes 22, 23 is managed by respective node manager components 29, 30.

In preferred embodiments of the present invention all entities in the framework illustrated in Figure 3 except the HostMasterWatchdog 19 are implemented as instances of Java™ classes, thereby providing an object oriented implementation. Details of such an implementation are described below. In the following description, programming language constructs such as classes, objects, interfaces and methods, are given their usual meaning within the Java programming language.

The HostMasterWatchdog process 19 regularly checks the state of the HostMaster process 20. If the HostMasterWatchdog process 19 finds a HostMaster process 20 in an invalid state, or fails to find a HostMaster process, it attempts to shutdown any existing HostMaster process 19 and starts a new HostMaster process. If this is initially unsuccessful, repeated retries are attempted until a user manually exits the HostMasterWatchdog 19. The HostMasterWatchdog process 19 is typically written in a conventional high level programming language such as C, and a HostMasterWatchdog implementation will therefore need to be written for each different platform on which the framework is to operate. However, in preferred embodiments of the invention, the other components are all implemented as instances of Java classes, providing portability between platforms.

Figure 4 illustrates the classes used to implement other entities illustrated in Figure 3. Individual classes are illustrated in further detail in Figures 4A to 4V.

In addition to the entities illustrated in Figure 3, each of the processes 21, 22 maintains a component registry by creating and maintaining an instance of a DefaultComponentRegistry class 31 (Figure 4A), which implements a ComponentRegistry interface 32 (Figure 4B). This registry allows objects to be added to the registry using a registerObject(Object C) method specified in the ComponentRegistry interface 32. Objects added to the registry can be monitored and

administered using methods provided by the DefaultComponentRegistry class 31, and therefore such components can receive notifications of state changes in other components in the registry.

Each of the entities illustrated in Figure 3 is represented by a corresponding class shown in Figures 4. The HostMaster process is represented by an instance of a HostMaster class 33 (Figure 4C). The ProcessAgent process is represented by an instance of a ProcessAgent class 34 (Figure 4D), which in turn references an instance of a NodeManager class 35 (Figure 4E) by means of a private variable mNodeManager within the ProcessAgent class 34. The NodeManager class 35 contains details of individual nodes. Classes used to represent such nodes will be described in further detail below.

Appropriate instances of the HostMaster 33 and ProcessAgent 34 classes set out above are registered with an instance of the DefaultComponentRegistry class 31. It can be seen that both HostMaster class 33 and ProcessAgent class 34 implement a ComponentAware interface 36 (Figure 4F). This specifies a single method GetComponentClass() which all implementing classes must provide. Each class defines this method so as to locate an appropriate corresponding component class. In the case of the HostMaster class 33, a HostMasterComponent class 37 (Figure 4G) is located by the GetComponentClass() method. In the case of the ProcessAgent class 34, the GetComponentClasss() method locates the ProcessAgentComponent class 38 (Figure 4H).

The component classes provide a convenient manner for encapsulating administrative information separately from functionality. When an object implementing the ComponentAware interface 36 is registered with the instance of the DefaultComponentRegistry class 31, the GetComponentClass() method is used to locate the appropriate component class, and this component class is then instantiated and registered with the instance of DefaultComponentRegistry. These component classes in turn links with the respective object which is being registered.

It should be noted that all entities to be registered with the DefaultComponentRegistry will not necessarily implement the ComponentAware interface 36. If the DefaultComponentRegistry object determines that an object being registered does not implement this interface, an instance of a GenericComponent class 39 (Figure 4I) is used to register the object with the DefaultComponentRegistry object. This class simply allows the object being registered to be located, but provides no administrative functions.

Figure 4 illustrates the hierarchy of classes used to represent components. The highest level of the hierarchy comprises a Component interface 40 (Figure 4J) which is implemented by a ComponentBase class 41 (Figure 4K). The ComponentBase class 41 has two subclasses, a G2Component class 42 (Figure 4L) and a AdminComponent class 43 (Figure 4M). The AdminComponent class 43 acts as a superclass for three classes which are used to implement administrative functions. A CascadingAgentComponent class 44 (Figure 4N) is used for communication between an instance of the HostMaster class 33 and an instance of the ProcessAgent class 34. An HTMLAdaptorComponent class 45 (Figure 4O) is used by the HostMasterComponent 37. An RMIConnectorServerComponent 46 (Figure 4P) is used by the ProcessAgentComponent 38 for communication via Remote Method Invocation (RMI) as provided by the Java programming language. The G2Component class 42 acts as a superclass for components described above. Specifically, the GenericComponent class 39, the ProcessAgentComponent class 38 and the HostMasterComponent class 37 are all subclasses of the G2Component class 42. Additionally, the G2Component class 42 is a superclass for a ServiceComponent class 47 (Figure 4Q), and a NodeComponent class 48 (Figure 4R).

It has been mentioned above that an instance of the NodeManager class 35 is referenced by a private variable in instances the ProcessAgent class 34. Both the ProcessAgent class 34 and the NodeManager class 35 both implement a NodeAdmin interface 49 (Figure 4S).

A number of other classes are also illustrated in Figure 4. A G2Runtime class 50 (Figure 4T) is used as a master class for the runtime of all classes illustrated in Figure 4.

A ComponentAddress class 51 (Figure 4U) is used to address the components of the architecture of Figure 3. An exception AdminException 52 is used for exceptions thrown by subclasses of the AdminComponent class 43.

An MBeanConfig class 53 (Figure 4V) is used to represent configuration information of an MBean™. The components implemented as subclasses of the ComponentBase class 41 all inherit a private variable of type Object which represents an MBean used to implement that component. The MbeanConfig class 53 represents the configuration of such MBean objects.

As described above, the HostMaster process 20 (Figure 3) is represented by a HostMaster object, and registered using a HostMasterComponent object. The HostMasterComponent allows the HostMaster to be notified of ProcessAgent creation and removal. This is described in further detail below.

The ProcessAgent processes 27, 28 (Figure 3) are represented by respective ProcessAgent objects which allow administration of nodes within that process, via appropriate NodeManager object(s).

As shown in Figure 3, processes contain nodes, which are managed by a NodeManager object. Nodes in turn comprise one or more services. A node is responsible for its contained services, and does not directly deliver any user functionality. Each node is represented by a suitable object, and Figure 5 shows an overview of appropriate classes. The individual classes are shown in Further detail in Figures 5A to 5J.

The NodeManager class 35 instantiates and uses an instance of a NodeFactory class 54 (Figure 5A) which creates instances of a MasterNode class 55 (Figure 5B), and these instances are used to represent the nodes 23, 24, 25, 26 (Figure 3). It should be noted that the MasterNode class 55 implements a Node interface 56 (Figure 5C) which specifies methods required to implement basic functionality. The NodeManager class instantiates instances of a NodeNotFoundException 35a, when it is requested to act upon a MasterNode object of which it is not aware.

Each MasterNode object has associated with it one or more router threads which are used to route messages within the node. These threads are represented by a private variable `mRouterThreadPool` in the MasterNode class 55, which is an instance of a class (not shown) implementing a Pool interface 57 (Figure 5D). The class implementing the Pool interface contains objects represented by instances of a RouterThread class 58 (Figure 5E), which represent the router threads.

Each node has associated with it a cache represented by an instance of a ReplicatedCache class 59 (Figure 5F) which is a subclass of a TimeStampedCache class 60 (Figure 5G), which in turn implements a Cache interface 61 (Figure 5H). The purpose of the ReplicatedCache object is to store state information related to each service contained in that node, and this information can be used to restart services in the event of a system crash. Furthermore, the cache is replicated between multiple nodes, such that if one node crashes, the node and its associated services can be restarted using cache data from the node containing the copy of the cache data. The ReplicatedCache class 59 provides methods to ensure synchronisation of data in the copy cache.

Each node also comprises an addressable registry, represented by an instance of the ComponentRegistry class 62 (Figure 5I). This class allows services within a node to be registered with the node, and also provides look up functions which can be used to locate services present within the node.

A node represented by a MasterNode object can have one of eight possible states, each of which is represented by a corresponding class. The private variable `mState` indicates an object indicating this state information. Each of the eight classes which may be stored in the `mState` variable implements the Node interface 56 which specifies methods relevant to transitions between states.

The state information is in fact represented by instances of classes which are subclasses of an AbstractNode class 63 (Figure 5J) which itself implements the node interface 56.

A class CreatedNode 64 represents a node which has been created but not yet initialised. A class InitializedNode 65 is used to represent a node following initialisation. A class

StoppedNode 66 represents an initialised node which is not currently running. When a node is starting, its state is represented by a class StartingNode 67. A class RunningNode 68 represents a node which has been started and is running. A node which is about to stop running is represented by a StoppingNode class 69. Classes FailedNode 70 and RecoveringNode 71 are used to represent node failure. Each of these classes specifies a constructor method taking as its sole parameter the MasterObject to which it belongs, and the also provides five methods specified in the Node interface 56 – init(), start(), stop(), destroy(), recover(). Each class representing state therefore provides a number of methods to allow state change. For each state, only a single method will not throw an InvalidStateException when called.

Transitions between the states represented by the objects set out above are now described with reference to Figure 6. Upon construction of a MasterNode object by instantiation of the MasterNode 55 by the NodeFactory 54, the mState variable points to an instance of the CreatedNode class 64.

Calling the init() method provided in the CreatedNode class creates an instance of InitializedNode 65, which is referenced by the MasterNode to represent state. When initialisation is complete an instance of the StoppedNode class 66 is created by the InitializedNode instance. Calling the start() method provided by the StoppedNode class 66 causes an instance of the StartingNode class 67 to be created which is then referenced by the MasterNode to represent state. The StartingNode class then creates an instance of the RunningNode class 68. The only method which provided by the RunningNode class 68 which does not throw the InvalidStateException is the stop() method. When this is called an instance of the StoppingNode class 69 is created, and subsequently a further instance of the StoppedNode class 66 is created to represent state.

If processing causes an uncaught exception to be thrown, an instance of the FailedNode class 70 is created. The only method which can then be validly used is recover() which creates an instance of the class RecoveringNode 71, before creating an instance of the class StoppedNode 66.

Referring back to Figure 5, it can be seen that the MasterNode class 55 implements a NodeStateSource interface 72 which allows instances of a NodeStateListener class 73 to be registered as listeners with MasterNode objects.

The preceding description has been concerned with the implementation of Nodes. It has been described that nodes contain services, and appropriate services are now described. Services are managed by a node, and receive and process messages which are sent to them. The class hierarchy used to implement services is now described with reference to Figure 7. Further details of the classes shown in Figure 7 are shown in Figures 7A to 7F.

Each service is implemented as an object which is an instance of a BaseService class 74 (Figure 7A), or as an instance of a subclass thereof. The BaseService class 74 implements an Addressable interface 75, which specifies a single getAddress method. In the case of the BaseService class 74, it can be seen that a service's address is stored as a private variable of type ComponentAddress (shown in Figure 4)

The BaseService class 74 additionally implements a HandlerContext interface 76 (Figure 7B) which specifies features which allow service handlers represented by instances of the BaseServiceHandler class 77 (Figure 7C) to be provided within a service to provide functionality. Handlers receive messages addressed to a service, and act upon those messages. Thus, handlers provide the functionality of a service. Each service will have a predetermined handler class which provides handlers, and this class will be a subclass of the BaseServiceHandler class 77. Each service may have a plurality of instances of the handler class, thus allowing a service to process messages from a plurality of users concurrently.

It can be seen that the BaseServiceHandler class 77 implements a MessageHandler interface 78 which allows handling of messages in a unified manner by specifying a single handleMessage() method. The BaseServiceHandler class 77 also implements a ResultListener interface 79 which specifies a single method deliverResult.

Instances of the BaseService class 74 use instances of a MessageDispatcher class 80 (Figure 7D) to dispatch messages within and between services. Both the BaseService

class 74 and the BaseServiceHandler class 77 implement a MessageRouter interface 81, which allows routing of messages between services using a route() method. This interface is used by the MessageDispatcher class 80. Routing of messages using these classes is described in further detail below.

It can be seen that Figure 7 additionally illustrates a ServiceFactory class 82 (Figure 7E) which is used to create services, and a Service State class 83 (Figure 7F) which is used to denote the state of an instance of the BaseService class 74 or one of its subclasses. Figure 7 also illustrates a DebugHandler class 84, which is a service handler used for debugging purposes during development.

Figure 8 illustrates a hierarchy of classes used to implement messages which can be passed between services to be handled by appropriate service handlers. These classes are shown in further detail in Figures 8A to 8K.

Referring to Figure 8, it can be seen that the hierarchy used to represent messages is headed by a Message interface 85 (Figure 8A) which is implemented by an abstract class AbstractMessage 86 (Figure 8B). A ComandSequenceMessage class 87 (Figure 8C) extends the AbstractMessage class 86. An instance of the CommandSequenceMessage class 87 represents messages which are to be transmitted between services. The CommandSequenceMessage class has four subclasses representing messages used in preferred embodiments of the present invention. A UserRequest class 88 (Figure 8D) is used to represent messages generated by a user. A UserResponse class 89 (Figure 8E) is used to represent responses to messages represented by instances of the UserRequest class 88. An ApplicationRequest class 90 (Figure 8F) is used to represent requests made using Application Programmers' Interface (API) calls, and an ApplicationResponse class 91 is used to represent response to messages represented by the ApplicationRequest class 90.

Each of the message types described above inherits core features required to implement messaging from the CommandSequenceMessage class 87 (Figure 8C). A private mSequence variable within the CommandSequenceMessage class 87 references an instance of a CommandSequence class 92 (Figure 8H). The CommandSequence class

includes an array of pairs which comprise a command, and a target (that is a service to which that command is targeted). Each command target pair is represented by an instance of a `CommandTargetPair` class 93 (Figure 8I). It can be seen that `CommandSequence` acts as a superclass for a `UserRequestCommandSequence` class 94, a `UserResponseCommandSequence` class 95, an `ApplicationRequestCommandSequence` class 96 and an `ApplicationResponseCommandSequence` class 97. Each of these subclasses of the `CommandSequence` class 92 represents a command sequence appropriate to the corresponding subclass of the `CommandSequenceMessage` class 87.

The state of any message represented by a subclass of the `AbstractMessage` class 86 is represented by a private `mState` variable which references an instance of a `MessageState` class 97 (Figure 8J).

The `CommandSequenceMessage` class 87 also has as a subclass a `SingleCommandMessage` class 98 (Figure 8K) which represents a message having a single command and a single target. The `SingleCommandMessage` class in turn has a `SystemRequest` class 99a and a `ServiceRequest` class 99b as subclasses.

Having described the various classes used to implement that which is illustrated in Figure 3, and having also described the classes used to implement messaging, the manner in which the components of Figure 3 operate together and the manner in which messaging is achieved is now described with reference to Figures 9 to 13. In the following description, objects are denoted by the reference numerals followed by a prime symbol ('). Where appropriate, reference numerals associated with the class of which an object is an instance are used.

Referring first to Figure 9, creation of `ProcessAgents` in the architecture of Figure 3 is illustrated. As described above a `HostMasterWatchDog` process 19 creates a `HostMaster` object 33'. The `HostMaster` object has associated with it a `G2Runtime` object 50' which can be used to locate a `DefaultComponentRegistry` object 31' using a `GetComponentRegistry()` method. The `HostMaster` object 33' is then registered within the `DefaultComponentRegistry` object 31' this in turn creates a `HostMasterComponent` object 37'.

The HostMaster uses configuration data to determine how many ProcessAgent objects it needs to create. For each ProcessAgent object which is to be created, a CascadingAgentComponent 44' is created, and this component in turn communicates with a respective ProcessAgent object 34' which is created within its own Java Virtual Machine. The main () method of the ProcessAgent object 34' is configured to carry out the work of the Process Agent. The ProcessAgent object 34' has an associated G2Runtime object 50'' which is used to locate a DefaultComponentRegistry object 31''. This in turn creates a ProcessAgentComponent object 38'.

The Component objects created by instances of the DefaultComponentRegistry class as described above allow administration of both HostMaster and ProcessAgent objects. As described above both the HostMaster object 33' and the ProcessAgent object 34' have associated DefaultComponentRegistry objects 31' and 31'', and both these DefaultComponentRegistry objects store details of both the HostMasterComponent object 37', and the HostMasterComponent object 38'.

The CascadingAgentComponent object 44' and the ProcessAgent object 34' communicate using RMI, and both have RMI clients which are used for this purpose. As described above, a HostMaster object is responsible for management of its associated ProcessAgent objects. This is achieved by communication using the component registries described above. Furthermore, a ProcessAgent object may generate "Heart beat" events at predetermined intervals, and the HostMaster object may listen for such events. If such an event is not received at the expected time, or within a predetermined time after the expected time, the HostMaster object then assumes that the ProcessAgent object has died, and accordingly performs recovery by instantiating a new ProcessAgent object. Details of the death of the earlier ProcessAgent object and creation of the new ProcessAgent object are recorded in the appropriate DefaultComponent registry objects.

Figure 10 illustrates the process of node creation within a particular process. The ProcessAgent object 34' locates its NodeManager object 35' (referenced by the mNodeManager private variable). It then uses a public createNode method within the

NodeManager object 35' to create a node. The createNode method takes as a parameter a node type, which is represented by a NodeType object which is an index and textual description of the node to be created. The NodeManager object 35' creates a NodeFactory object 54' using its constructor which takes no parameters. The createNode method provided by the NodeFactory object 54' is then used to create the desired node. The NodeFactory object 54' creates a MasterNode object 55' by calling its constructor method which takes no parameters.

The MasterNode object 55' in turn instantiates a CreatedNode object 64' which is referenced by an mState parameter within the MasterNode object 55'. An init() method provided by the CreatedNode object 64' is then called by MasterNode to initialise the MasterNode object 55', and the transitions between states illustrated in Figure 6 can then occur as necessary, and the state of the node is set using a setState method provided by the MasterNode object 55'. The CreatedNode object 64 in due course uses a ServiceFactory object 82' to create the services required by the MasterNode object 64'.

It can be seen from Figure 4E that the NodeManager class includes a private Vector variable mynodes which stores details of all nodes managed by a NodeManager object. Figure 10 schematically illustrates that each time a MasterNode object is created, it is added to this vector variable.

Figures 11A and 11B show how messages are transmitted from a first service 74' to a second service 74'' in accordance with an embodiment of the present invention. Referring to Figure 11A it can be seen that the first service 74' comprises a plurality of service handlers 77' which provide the functionality of the first service as described above. The first service 74' additionally comprises an in queue IN1 into which incoming messages are placed, and an out queue OUT1 into which outgoing messages are placed prior to transmission. The second service 74'' comprises a plurality of service handlers 77'', an in queue IN2 and an out queue OUT2. Both the first service 74' and the second service 74'' have associated with them unique addresses which are used for directing messages to them.

Figure 11B illustrates how messaging is effected using the structure illustrated in Figure 11A. Figure 11B illustrates routing of a message from one of the service handlers 77' of the first service 74' to one of the service handlers 77'' of the second service 74''. At step S1, a service handler 77' of the first service 74' executes the command specified in the message. When this execution is complete, the service handler 77' interrogates the message to determine an address mask of the next service to which it is to be sent, and modifies the message to show this address mask (step S2). At step S3, the message is sent to the out queue OUT1 of the first service 74'. A MessageDispatcher object associated with the out queue OUT1 then copies the message to its wait queue (step S4), and interrogates the message to obtain an address mask indicating the message's intended destination (step S5). At step S6, a ComponentRegistry object is used to locate all services within the current operating system process which match the message's address mask.

At step S7, a check is made to determine whether or not any services have been located. If it is determined that a single appropriate service has been located (e.g. the second service 74'' of Figure 11B), the message is forwarded to the in queue of that service (step S8), an appropriate service handler is located within that service (step S9), and the message is forwarded to the located service handler.

If step S7 determines that no suitable service can be found within the current process, the message is amended to show that it must be directed to a messaging service which can carry out inter-process messaging (step S11). An appropriate messaging service is then found (step S12), the message is directed to the in queue of the messaging service (step S13) and the messaging service then carries out inter-process communication (step S14).

If step S7 locates N suitable services, where N is greater than one, a counter variable, m, is initialised to zero, (step S15), and then a sequence of operations are carried out N times by the action of a loop established by steps S16 and S21. Each loop will produce a clone of the message (using standard Java functionality) at step S17 and this is sent to one of the located services at step 18, A suitable service handler is then located within each service (step S19) and appropriate processing is carried out (step S20).

Figures 12A to 12D illustrate messaging in further detail.

Referring to Figure 12A, it can be seen that a `CommandExecutor` object 77b' (referenced by an appropriate variable in an appropriate handler class) uses an `execute` method provided by a command stored in a `Command` variable within a message 88a'. In the example illustrated in Figure 12A the command is a `ReceiveUserRequestCommand`, which in turn is handled by a connection handler object 77c'. It is at this stage that service and message specific processing takes place. Assuming that the command is correctly executed, the `CommandExecutor` object 77b' uses a `commandSucceeded` method provided by the message 88a'. It can be seen that this method is provided by the `AbstractMessage` class (Figure 8B), and all messages are instances of subclasses of the `AbstractMessage` class. In the case of a `CommandSequenceMessage`, the `commandSucceeded` method will cause a `CommandTargetPair` variable `m_ProcessingDelivery` to be set to the next command, and the next target address. When this is complete, a `SendMessage` method provided by a `ConnectionService` object 74a' is used to direct the message to the service's out queue (represented by an `mOutQueue` variable in the `BaseService` class), using a `put` method. Thus Figure 12A illustrates the steps required to place a message in a service's out queue.

Referring to Figures 8B and 8C it can be seen that a `CommandSequenceMessage` object comprises a command sequence (which is an instance of a `CommandSequence` object, Figure 8H), and three `CommandTargetPair` objects. A first `CommandTargetPair` is set in the case of success as described above using the `commandSucceeded` method. A second `CommandTargetPair` is set in case of failure by a `commandFailed` method, and a third `CommandTargetPair` is used where communication is required between processes, as is described below with reference to Figure 12D. State information within a message object is used to determine which of the `CommandTargetPair` objects should be used at a particular time.

Figure 12B illustrates steps which are carried out after a message has been placed in a service's out queue. The out queue of the service is provided with its own

MessageDispatcher object 80', which implements the Thread interface provided by the Java programming language. A get() method provided by a WaitQueue is called by a MessageDispatcher object 80'. The get() method transfers the message to a queue represented by an mWaitQueue object 80a' within the Message dispatcher object 80'. The MessageDispatcher object 80' then locates an appropriate MasterNode object 55', and routes the message to that node.

On receiving the message, the MasterNode object 55' locates a RouterThread object 58' from its pool of router threads (represented by an instance of a class implementing the Pool interface illustrated in Figure 5D.), using a getRouterThread method. Having located a suitable RouterThread object, a route method provided by a RouterThread object 58' is used to route the message to the correct service. This involves calling the getAddressMask method provided by the Message object to obtain an address mask, and then using a ComponentRegistry object 31' to locate all services matching the obtained address mask by calling a findLocalAddressables method provided by the ComponentRegistry object 31'. Assuming that at least one suitable service is found, a getRouter() method is called on the Service object 74b'. A route() method provided by the Service object 74b' is then used to route the message to out queue 74c' of the Service object 74b'.

If no suitable services are located by the findLocalAddressables method, the processing illustrated in Figure 12D is carried out, as described in further detail below.

If more than one service is found within a node which matches the specified address mask, then the message is cloned and sent to the in queues of all services in the manner described above. Furthermore, it should be noted that as soon as the message is dispatched to the RouterThread object 58' the work of the MessageDispatcher object is finished, thus operation of the MessageDispatcher object 80' is decoupled from operation of the RouterThread object 58'.

Figure 12C illustrates processing when a message has reached a service's in queue. A MessageDispatcher object 80b' is associated with the in queue and listens for message arrivals. When a message arrives, the MessageDispatcher object 80b' uses a get()

method to copy the message to its WaitQueue object 80c'. A route method is then used to forward the message to a MessageToHandlerForwarder object 80d'. The MessageToHandlerForwarder object 80d' uses a getObject method provided by a ServiceHandlerPool object 80e', to provide a Handler object 77d'. The MessageToHandlerForwarder object 80d' then call a handleMessage method (specified in the MessageHandler interface) to cause the message to be handled. In due course the Handler uses a setNextCommand method provided by a CommandExecutor object 77b' to cause the command sequence to be suitably updated.

Figure 12D illustrates the processing shown in Figure 12B, modified where it is the case that the RouterThread object 58' determines that there are no services registered with the ComponentRegistry 31' which match the specified address mask. In this case, a leaveProcess method is used to update the internal state of the message object 88a', and the RouterThread then seeks to locate a service which is responsible for inter-process communication by using a getAddressMask method and a findLocalAddressables method, which will return an appropriate service 74d'. Having located an appropriate service the message is sent to that service in the manner described above with reference to Figure 12B. On receipt of the message in the In queue of the service, the service handles the message (as described with reference to Figure 12D), and inter-process communication is thus achieved.

From the preceding description, it can be seen that message routing is carried out as specified within appropriate variables of a given message. There is no central control of message routing, but instead control is distributed between objects which handle messages.

For example, referring to Figure 13, there is illustrated a system of five message handlers denoted A, B, C, D and D'. The message handlers D and D' are both instances of a common message handler.

A message Msg1 is sent from handler A, to handler B, and subsequently from handler B to handler C. Handler C processes Msg1 in a predetermined manner, and generates two messages, Msg2 and Msg3. Each of these messages is then routed independently. The

message Msg2 is routed to handler D while the message Msg3 is routed to the handler D'.

Each handler processes the message, and routes the message onwards to the next service specified therein. Distributing control of message routing in this way provides a scalable system, and removes the need for any centralised control of messaging. This is particularly beneficial where a handler creates a plurality of messages from a single received message (as in the case of handler C in Figure 13).

Referring back to Figure 2, having described a framework for operation of the server 13, operation of the webserver 12 and server 13 to provide composite user interfaces is now described.

Figure 14 illustrates the logical architecture of the web server 12 and the server 13. It can be seen that the server 13 operates using a process framework as illustrated in Figure 3, and comprises a Host Master Watchdog process 100, a Host Master process 101 and a process 102.

The process 102 comprises a process Agent 104 and a Node Manager 105 which perform functions as described above. The process 102 comprises three nodes 106, 107, 108. A client adaptor (CA) node 106 is responsible for interaction between the server 13 and the composite application(s). An application adapter (AA) node 107 is responsible for interactions between the server 13 and the appropriate source applications via a source application interface 109. An ALF (Authentication Lockdown Framework) node 108 is responsible for authentication and security functionality provided by the server 13. Security information is stored in a database 110 which is accessed via a conventional Java Database Connectivity (JDBC) interface 111. The function of each of these nodes is described in further detail below.

The server 13 is connected to an administration terminal 112 via a link 113. The administration terminal 112 allows administrators of the server 13 to perform various administration functions, as described below.

Configuration data for the nodes 106, 107 is obtained from an Integration Data Manager (IDM), which comprises a database 115 accessed via an appropriate interface (such as JDBC interface 116). The form of the configuration data, and the manner in which it is used is described in further detail below. It should be noted that the database 110 and the database 115 together make up the configuration data 18 of Figure 2.

The Webserver 12 communicates with the server 13 via a connection 14. The webserver runs a Java™ servlet 117, and this servlet 117 is responsible for communication between the webserver 12 and the server 13. In many situations is necessary to impose restrictions on access to composite applications or parts of composite applications managed by the server. Such restrictions can be conveniently implemented by allocating user names and passwords to users, and associating appropriate access privileges with such usernames and passwords. The servlet 117 implements any necessary security policies, by cooperating with the CA node 106. The functionality provided by the servlet 117 is referred to as a User Agent (UA) service.

The CA node 106 comprises four services. A Messaging layer (ML) service 118 provides Java Messaging service (JMS) functionality for communication between objects within different processes. The JMS is such that messaging between processes can be carried out in a unified manner, regardless of whether processes are located on the same or different hosts.

A Communications Layer (CL) service 119 provides connections between the CA node and other processes. A Data Transformation (DT) service 120 provides means to transform data from one form to another. Such a service can be required both when receiving user data from a composite application which is to be forwarded to a source application, and when receiving data from one or more source applications which is to be output to a composite application. A User Experience Manager (UXM) service 121 is responsible for determining the action necessary when requests are received from the webserver 12, and when data is received from a source application.

The AA node comprises three services an ML service 122, a CL service 123 and a DT service 124. Each of these services provides the same functionality as the equivalent service of the CA node 106.

The ALF node 108 again comprises an ML service 125 and an ALF service 126. The ALF service 126 is responsible for security features as is described in further detail below.

The use of the CA node 106 and the AA node 107 to produce and manage composite applications is now described with reference to Figure 15. A user operates a web browser 126 to access HTML documents provided by the web server 13 which make up a composite user interface. A user issues a request via the web browser 126. This request is forwarded to the webserver 12 using the Hyper Text Transfer Protocol (HTTP) operating in a conventional manner. The servlet 117 within the web server 12 operates a UA service 127. The UA service 127 authenticates the request in accordance with predetermined security policies (defined as described below). If the authentication is successful, an IncomingUserRequest message is created using parameters received by the UA (e.g. from the HTTP request). This IncomingUserRequest message is then sent by the UA service 127 to the server 13, having a target address of the CL service 119 of the CA node 106. As described above, this message is transmitted using the JMS protocol. The CL service 119 receives the IncomingUserRequest message, and having received the message, directs it to the DT service 120 of the CA node 106. The DT service 120 performs any necessary transformation of data contained in the IncomingUserRequest message, and directs the message onwards to the UXM service 121 of the CA node 106. The UXM service 121 processes the message and generates one or more OutgoingUserRequest messages, destined for the appropriate source application(s) 128. These messages are first sent to the DT service 124 of the AA node 107 using, for example, the JMS protocol.

The DT service 124 of the AA node 107 performs any necessary transformation, and forwards the message to the CL service 123, which in turn generates messages to request appropriate data from the source application(s) 128 in an appropriate format.

These messages are then forward to the Source Application(s) using, for example, the JMS protocol.

The Source application(s) 128 process received messages and in turn produce data in response to the received request. This data is received by the AA node 107, and passed to the CL service 123. The data is then forwarded to the DT service 124 which processes the data to extract one or more page objects and to create an IncomingUserResponse message which is forwarded to the UXM service 121 of the CA node 106.

The UXM service 121 processes the IncomingUserResponse message, and determines whether or not it has received all data necessary to create a composite page (determined by its configuration, which is described below). When all data required is received, an OutgoingUserResponse message is created which is forwarded to the DT service 120 where any necessary transformation is carried out. Having performed any necessary transformation the OutgoingUserResponse message is forwarded to the CL service 119 and then onwards to the ML service. The ML service transmits the composite page to the UA 127 using the JMS protocol. The UA then displays the page to the user via the web browser 126, assuming that any necessary validation is successfully carried out by the UA 127.

In the description presented above, a distinction is made between IncomingUserRequest messages and OutgoingUserRequest messages. A similar distinction is made between IncomingUserReponse messages and OutgoingUserResponse messages. These distinctions are made to aid understandability, however it should be appreciated that in some embodiments of the invention both IncomingUserRequest messages and OutgoingUserRequest messages are represented by a single UserRequest message type. Similarly both IncomingUserReponse messages and OutgoingUserResponse messages may be represented by a single UserResponse message type.

It will be appreciated that the processing described above can be used to handle a wide range of different composite applications, however it will also be appreciated that the processing required will vary considerably, and therefore an effective means of

configuring the services illustrated in Figures 14 and 15 is required. This is achieved by accessing the IDM database 115. IDM data is stored in a hierarchical manner as will be described in further detail below.

At startup, the web server 12 and the server 13 need to be provided with data which allows access to the IDM database, and indicates where in that database the appropriate data is located. This is achieved by providing each with appropriate invocation parameters which indicate where the configuration data can be located.

The servlet 117 of the webserver 12 (referred to as “the listener”) is provided with invocation parameters as illustrated in Figure 16.

A `g2jms.config` parameter specifies a file containing JMS configuration properties for the servlet 117. A `g2config.home` parameter specifies a directory containing an IDM properties file, which provides information as to how the IDM database 115 is accessed. A `g2listener` parameter specifies a node in the hierarchical IDM data structure where configuration information for the UA service 127 is located. A `g2config.instance` parameter species configuration data for the composite application which the webserver 12 is to provide to a user. A `g2jms.config` parameter specifies configuration for the JMS service of the UA 127. A `g2tracer.conf` parameter specifies logger configuration.

The nodes and services of the server 13 (collectively referred to as “the composer”) are provided with invocation parameters are illustrated in Figure 17.

The `g2config.home`, `g2jms.config` and `g2config.instance` parameters perform functions as described with reference to Figure 16. The `g2config.webhost` parameter specifies an IP address for the web server on which the UA service 127 is implemented by means of the servlet 117. The `g2config.baseipport` parameter specifies a port number used by an administration console which is used in connection with the composer as described below. The `g2basetracer.conf` parameter is used to specify a log file for each process operating within the composer. Given that every composer comprises at least a HostMaster process and a process containing AA and CA nodes, at least two log files must be specified, in addition to a log file for the web server.

The parameters described above with reference to Figures 16 and 17 together allow the listener and the composer to obtain configuration data from the IDM. Additional invocation parameters are illustrated in Figure 18. These specify page tags which are used in HTTP requests passed to the listener by the composer.

APPLICATION_CLASS specifies a name used for the composite application, and is used by the composer to determine any transformation that the CA.DT may be required to perform. UXM_TREEID and UXM_NODEID specify where in the configuration the UXM can find configuration information pertinent to the specific composite page. UXM_MODIFIED allows the UXM to perform no processing on some pages. If UXM_MODIFIED is set to true, this is an indication that processing is required in accordance with the configuration data stored in the IDM (described below). If UXM_MODIFIED is set to false, this means that no processing is required and the UXM simply forwards the message to the AA node.

Additional tags are used for security purposes. A g2authorization tag is used to indicate that the listener should attempt authentication. usr and pwd tags are respectively used to represent user name and password data when this is required for authentication.

The invocation parameters described above allow the composer and the listener to locate the IDM and to locate suitable data within the IDM. These parameters are conveniently implemented as variables of appropriate Java classes. They can be configured by means of a system.properties file, an example of which is illustrated in Figure 19. Each line of the file corresponds to a Java property to be set. It can be seen that all entries of the file are prefixed by either 'hm' or 'pa'. This ensures that each entry is picked up by the correct process. Entries prefixed by 'hm' are picked up by the HostMaster process, and it can be seen that these correspond to the composer invocation parameters described above. Entries prefixed by 'pa' are additional properties to be set in each process set up by the HostMaster process. These typically include configuration for various aspects of the Java Virtual Machine (JVM) such as garbage collection.

It can be seen that in Figure 19, all entries which begin 'pa' are of the form 'pa.x'. This indicates that the property should be set on all processes started by the HostMaster. However, in some embodiments of the invention entries may be prefixed by 'pa.N' where N is an integer. In this case, that entry refers only to process N.

It should be noted that the system.properties file is used in a hierarchical manner as follows. Any entry of the form 'pa.N' will override a 'pa.x' entry for process N. Any entry of the form 'pa.x' will override the 'hm' entry inherited from the HostMaster process. In addition to setting parameters by means of the system.properties file, it should be noted that configuration parameters can be specified at start up, for example via a command line interface. In such a case, if a property value is specified both in the system.properties file and on the command line, the system.properties file will determine the value.

Having described invocation parameters, and associated tags, the structure of IDM data stored in the database 115 is now described.

It should be noted that, as will become apparent below, every composite application has a unique application class name which is used to identify it. Every source application also has an unique application class name which is used for identification. Thus, if a composite application is made up of two source applications, its configuration will include three application class names.

A partial IDM structure for a configuration containing a single process containing an AA node and a CA node (as illustrated in Figure 14) is shown in Figure 20. The highest level of the hierarchy comprises a single entity CONFIG 129 which is the root of the IDM tree. This entity has six child entities. An entity INSTANCE 130 is concerned with configuration information for a particular composer configuration. An entity BASE_ALF_CONFIG 131 provides configuration information for an ALF node. Entities HOST 132, PROCESS 133, NODE 134, and SERVICE 135, each have children containing configuration information for the corresponding entities of Figure 14.

As described above, the g2instance parameter references a named node in the IDM that contains appropriate configuration data. This data will include details of the number of processes that should be used, the allocation of CA nodes, AA nodes and ALF nodes to those processes, and locations where data appropriate for the services contained within those nodes can be found.

In this case the g2instance invocation parameter is set such that:

g2instance=LAB1_INS

This indicates that the composer should locate the LAB1_INS entity 136 which is a child of the SINGLE_PROC_AA_CA_INSTANCE 130a, which is itself a child of the instance entity 130a to obtain configuration data. It can be seen that the LAB1_INS entity has suitable data provided by a g2config data group 137.

The entry alf in the data group 137 indicates that server uses the ALF, and provides means for locating the ALF database 110 (Figure 14). The entry host indicates that the server 13 comprises a single host referenced by that entry in the data group 137. It can be seen that all other entries are of the form "host.process1.". Each of these entries provides details of a respective node or service operating within the single process 102 (process1), on the server 13 of Figure 14. It can be seen that every node and service illustrated in Figure 14 has a corresponding entry in the data group 137, which provides details of its configuration.

Configuration of the CL service 123 of the AA node 107 (Figure 14) is now described. Figure 21 shows part of the hierarchy of Figure 20 in further detail, together with details of entities relevant to the configuration of the CL service 123 of the AA node 107. In addition to the entities described above, the hierarchy comprises a BASE_CL_SERVICE entity 138 which is a child of the SERVICE entity 135. The BASE_CL_SERVICE entity 138 has two child entities, a AA_CL_SERVICE entity 139 which represents the CL service 123 of the AA node 107, and an EXT_APPS entity 140 which contains details of relevant source applications. It should be noted that the hierarchical arrangement of the IDM allows child entities to inherit and override

properties of their parent entity. For example, default data for all services may be specified in the SERVICE entity 135. Some this data may be inherited and used by the BASE_CL_SERVICE entity 138, while other parts of the data may be replaced by more appropriate values specified in the BASE_CL_SERVICE entity 138.

It can be seen that the host.processxAA.CL entry in the data group 137 references a data group g2config 141 under a mysrcapps entity 142. The g2config data group 141 includes an entry for every source application with which the AA communicates, and each application is referenced by its application class name. Each entry will include a session_timeout parameter which indicates a time (e.g. in minutes) which a user can be inactive without a connection being terminated, and a replicate_sessions parameter which is a Boolean indicating whether or not session data should be copied between processes. Setting this parameter to TRUE provides fail over if session data is lost for any reason.

The source applications from which the composite application is created are each represented by a child entity of the EXT_APPS entity 140. In Figure 21, two source applications are illustrated. A first source application (having a class name "SrcApp1") is denoted by an entity SrcApp1 143 and a second application (having a class name "SrcApp2") is denoted by an entity SrcApp2 144.

Each of these entities has two data groups. The SrcApp1 entity 143 has a g2config data group 145 which is referenced by the data group 141 and which specifies data for the configuration of the AA, and a connection data group 146 which specifies communications protocols for communication with the SrcApp1 application. The entity SrcApp 2 144 has two similar data groups 147, 148.

The data stored in relation to each source application is now described. The g2config data group for each source application 145, 147 includes a number of parameters. An authorization parameter takes a Boolean value and indicates whether or not the source application requires authentication. A credential_type parameter is required when the authorization parameter is set to TRUE. The credential_type parameter is either set to USRPWD or DYNAMIC. When credential_type is set to USRPWD, a static username,

password combination is used for authentication. When credential type is set to DYNAMIC, a static username and dynamically created password is used for authentication. For example, a password can be dynamically created from a user name using a predetermined algorithm. A protocol parameter indicates a protocol to be used for communication with the source application. Suitable protocols may include HTTP, SOAP, MQSERIES and JDBC. Other parameters provide information as to how errors should be handled.

The connection data group associated with each source application 146, 148 stores information related to each connection. It can be seen from Figure 21 that the connection data groups are not directly referenced by the g2config data group 141, and are located by virtue of the fact that they are located below the same entity as the g2config data group 145, 146 for each application.

Each connection data group 146, 148 will include indications of a protocol to be used, a port number of the source application to which data should be sent, an identifier for the host on which the source application is operating, a proxy identifier (if appropriate), together with parameters related to session refresh and logout.

In addition to the parameters described above, the connection data group will include various other parameters appropriate to the protocol being used.

If SOAP is being used these parameters will include a file part of a URL to use for connection, a URI style for message encoding, a URI for the target object, and a URI indicating the intent of the SOAP request.

If JDBC is being used, these parameters will include a URL for JDBC database connection, a JDBC driver name, and parameters indicating both initial and maximum JDBC connection pool sizes.

Other protocols will require other parameters, and such parameters will be readily apparent to those of ordinary skill in the art.

In addition to the data set out above, various other configuration data may be used to configure the CL service of the AA. Such configuration data may either be stored in one of the data groups described above, or alternatively may be stored in an additional data group. Typically, such additional data groups are located within the same source application entity 143, 144 as the data groups described above, and accordingly the instance entity 136 is able to locate all necessary configuration data.

The CL service 119 of the CA node 106 (Figure 14) is configured in a similar manner, although it will be appreciated that in this case, data of a similar form will be required for the or each composite application instead of the source applications. In some embodiments of the IDM, the data group 141 includes an entry for each composite application in addition to entries for each source application, and these entries in turn identify appropriate entities located below the EXT_APPS entity 140.

Configuration data for the DT service 120 of the CA node 106, and for the DT service 124 of the AA node 107 is now described. Figure 22 is a tree diagram showing part of the hierarchy of Figure 20, together with details of entities pertinent to configuration of the DT services.

The service entity 135 has a child entity 149 BASE_DT_SERVICE which in turn has a child entity DT_SERVICE 150 which represents all DT service configuration data. An entity LAB1_DT 151 is a child entity of DT_SERVICE and represents DT service configuration for the application represented by the LAB1_INS entity 136.

It can be seen that the host.process1.AA.DT and host.process1.CA.DT entries in the data group 137 both reference a DT.service data group 152 which is located beneath the Lab1_DT entity 151. The DT.service datagroup contains a single entry which allows it to locate its parent entity.

A DT.Applications data group 153 is located under the same entity 151 as the DT.service data group 152. The DT.Applications data group 153 contains one entry for each composite application which must be handled by the DT service of the CA, and one entry for each source application which must be handled by the DT service of the

AA. In the example of Figure 22, it can be seen that the data group 153 includes an entry for a single composite application Lab1App, and a single source application firstbyte. Figure 22 shows data needed to configure the DT service of the CA to handle the composite application Lab1App.

It can be seen that the Lab1App entry in the DT.Applications data group 153 references a DT.AppLab1App data group 154 which contains data relevant to handling transformations to and from the composite application Lab1App. The DT.App.Lab1App data group 154 includes an entry indicating transformations required in response to an IncomingUserRequest, an entry indicating transformations required in response to an OutgoingUserRequest, and an entry indicating transformations required to convert data from the UXM in a form suitable for output to the composite application.

The UXM entry references a DT.App.Lab1App.UXM data group 155 which in turn contains entries which indicate how data should be transferred between UXM objects used in the UXM of the CA, and data which can be processed by the composite application. In the example of Figure 22, the data group 155 includes a NewObjectTransformations entry which references a data group 156 and a UIElementTempates which references a data group 157. This data allows new data to be added to data received from source applications to create a composite page. The UIElementTemplates data group 157 contains details of such new data, and the NewObjectTransforamtions data group indicates how this data should be included in a composite page.

It can be seen that the other entries in the DT.App.Lab1App data group reference other data groups 158, 159, 160 which contain data indicating how these transformations should be carried out.

Figure 23 shows part of the IDM data structure which holds data pertinent to the DT service of the AA node. It can be seen that the firstbyte entry in the DT.Applications data group 153 references a DT.App.firstbyte data group 161, which includes an entry for each data type which the DT service of the AA may be required to handle, that is IncomingUserRequest, OutgoingUserRequest, IncomingUserResponse,

OutgoingUserResponse, and UXM objects. It should be noted that although in the example of Figure 15 only OutgoingUserRequest and IncomingUserResponse messages are processed by the DT service of the AA. However, IncomingUserRequest and OutgoingUserResponse messages are included for the case where the composer simply passes messages between a composite application and a source application without carrying out any processing.

It can be seen that the entries in the DT.App.firstbyte datagroup 161 relating to the four message types reference appropriate data groups 162, 163, 164, which contain the necessary transformation information. The UXM entry references a DT.App.firstbyte.UXM data group 165, which includes four entries. An ErrorPageResolverData entry references a DT.App.firstbyte.UXM.ErrorPageResoverData data group 166 which contains means to recognise various error pages which may be generated by the firstbyte source application. A PageResolverData entry references a DT.App.firstbyte.UXM.PageResolverData data group 167 which contains data indicating how pages of the source applications are recognised, in order to identify necessary transformations. A PageExtractionConfig entry references a DT.App.firstbyte.UXM.PageExtractionConfig data group 168 which contains references to a plurality of Extensible Markup Language (XML) files, which indicate how each page recogniser by an entry in the PageResolverData data group should be transformed by the DT service before being passed to the UXM service of the AA.

Configuration data for configuring the UXM service 121 of the CA node 106 is now described with reference to Figure 24, which shows the relevant parts of the IDM data structure. The highest level of the illustrated tree structure is a BASE_UXM_SERVICE entity 169 which is a child of the SERVICE entity 135 (Figures 20 to 23). The BASE_UXM_SERVICE entity 169 has three child entities, a CA_UXM_SERVICE 170 which represents data relevant to the UXM service 121 of the CA node 106 (Figure 14), A UXMAActionLibrary entity 171 which stores UXM action data, and a UXMPredicateLibrary entity 172 which stores predicate data.

The CA_UXM_SERVICE entity 170 has a child entity named TreeRootNode for each composite application which the UXM is able to handle, and an additional entity 173 to deal with error conditions. A TreeRootNode entity 174 represents a first composite application CompApp1, while a second TreeRootNode entity 175 represents a second composite application. Each TreeRootNode entity 173, 174, 175 has a control data group which specifies a unique identifier. In the case of the data group 176 associated with the entity 173, this identifier is simply "error", while in the case of the data groups 177, 178 an appropriate identifier is set. The control data groups additionally include information which can be used to limit a particular node's usage. Under the TreeRootNode 174 which relates to CompApp1, there are two child TreeNodes which each represent different pages of the composite user interface for CompApp1. A first child TreeNode 178 and its associated data group represent a page customer page CustPg, while a second child TreeNode 179 and its associated data group represent an order page OrderPg. It can be seen that the TreeNode 178 in turn has three child TreeNode entities 180, 181, 182, each of which represent particular parts of the customer page.

It was indicated above that the identifier within the control data group for each composite application must be unique. It should be noted that within a particular application, all identifiers must be unique within that application.

Figure 25 shows an alternative view of the tree structure of Figure 24. It can be seen that in addition to the control data groups described above, All TreeRootNode entities 174, 178, 180, 181, 182 additionally include an integration data group 183 which specifies the actions to be taken, if any, when a composite application of the type represented by the TreeRootNode 174 is encountered. This information is determined by data stored in a library located under the UXMAActionLibrary entity 171.

Additionally, the TreeRootNode entities 178, 180 comprise a predicates data group which specifies conditions (or predicates) which need to be true for the actions specified in the integration data group to be taken. Each of these predicate data groups references a predicate library located under the UXMPredicateLibrary entity 172.

Figure 26 shows part of the UXM configuration for a composite application Lab2App represented by a TreeRootNode entity 183. It can be seen that the TreeRootNode entity 183 has a control data group 184 and an integration data group 185. A TreeNode 186 specifies UXM data for part of the Lab2App application, and has a control data group 187 and an integration data group 188. The integration data group specifies actions to be taken in response to particular events, and these actions are specified in terms of an appropriate data group. For example, an entry for an action deleteUXMObj.1 references a DeleteUXMObjectFromTagetConext.RemData data group 189 within an action library Lab2_Lib 190 for the Lab2App application. Similarly an aggregateNews.2 entry in the integration data group 188 references an AggregateNamedContexts.AggregateNews data group 191 in the Lab2_Lib library 190. In turn, this data group references two TreeNodes 192, 193 which are child entities of the TreeNode entity 186. The other entries in the integration data group 188 reference corresponding entities 194, 195 in the Lab2_Lib library 190. Operation of the UXM in accordance with the configuration data described above is set out in further detail below.

Configuration of the ML services of the CA node 106 and the AA node 107 (Figure 14) is now described. As indicated above, communication between processes is accomplished using JMS, and the JMS protocol is encapsulated within the ML services. Referring back to Figure 20 it can be seen that ML entries of the g2config data group 137 reference an appropriate g2config data group 196, located beneath an ML_SERVICE 197 which is itself located beneath a BASE_ML_SERVICE entity 198. For the most part, configuration of the ML is based upon properties files, not by data within the IDM. These files can be located using invocation parameters, as described above with reference to Figures 16 and 17. The ML is configured using properties files instead of IDM data because much of the configuration is required at start up, before the IDM has been accessed. In preferred embodiments of the invention, a common configuration is shared by the CA node 106 and the AA node 107. JMS makes use of the Java Naming and Directory Interface (JNDI) to obtain necessary information. The properties files must therefore contain details which allow connection to the JNDI. Additionally, the properties file must include parameters which specify all necessary configuration data for the JMS implementation used to put embodiments of the

invention into effect. In preferred embodiments of the present invention, the SunONE MQ3.0 JMS implementation or the OpenJMS 0.7.2 implementation are used. Configuration data required will be well known to those skilled in the art, and is not described in further detail here.

However, it should be noted that some configuration parameters used in preferred embodiments of the present invention provide beneficial results. For example, it should be noted that each process within the system can send broadcast messages using a broadcast topic. Each process will have a single in queue and a broker will manage these queues ensuring that any message is sent to the in queues of all relevant processes. Additionally, each process has a persistence queue, into which messages are delivered if the in queue is unable to accept messages for any reason. This queue provides failover for the system.

Having described configuration of all services of the CA node 106 and the AA node 107 (Figure 14), configuration of the ALF service 126 of the ALF node 108 is now described. ALF configuration data is stored in the hierarchical IDM data structure, as illustrated in Figure 27. The g2config data group 137 for the entity Lab1_INS 136 references a g2config data group 199 located beneath the BASE_ALF_CONFIG entity 131. The g2config data group 199 provides the information necessary to allow the ALF database 110 (Figure 14) to be located. It can be seen that this data group includes details of a url for the database, a driver to be used with the database (here the Oracle™ JDBC driver) and user name for use when logging onto the database 110. It will be appreciated that other data may be required in order to properly access the database 110 and this too will be stored in the g2config data group 199.

Each user who is authorised to use a composite application will have a user entity 200, 201 which is a child of the BASE_ALF_CONFIG entity 131. The user entity 200 has a G2 data group 202 which includes a user name used by that user when logging on to composite applications provided by the composer. Additionally, the user entity 200 has a SrcApp data group 203 which includes details which are needed to log on to a source application "SrcApp". A similar data group is provided for each source application

which a user is able to access. The user entity 301 has a G2 data group 204, and will also comprise one or more entries for appropriate source applications (not shown).

The BASE_ALF_CONFIG entity 131 also has a usersconfig data group 205 which specifies details relating to logging onto to the composite application, for example password format, and a frequency with which passwords must be changed. The values specified in the usersconfig data group 205 provide a default log on configuration, which is inherited by all user entities. However, this default configuration may be overridden by a usersconfig data group located below an appropriate user entity. It should be noted that similar information relating to each source application is stored as part of the configuration of the CL service within the AA node.

The BASE_ALF_CONFIG entity 131 has an additional child entity CONSOLE_CONFIG 206. This entity comprises a plurality of data groups (not shown) which state how configuration data input to the system via the administration terminal 112 (Figure 14) is converted to ALF commands. Commands are typically received from the web browser based administration terminal 112 as HTTP requests, which are mapped to appropriate ALF commands. The administration terminal 112 is used to add users and perform other necessary administration functions related to the ALF.

Referring back to Figure 20, it can be seen that the host.process1.ALFNODE entry in the g2config data group 137 references a g2config data group 206 located beneath BASE_ALF_ML_NODE 207. The host.process1.ALFNODE.ALF entry in the g2config data group 137 references a g2config data group 208 located beneath a BASE_ALF_SERVICE entity 209. The g2config data group 208 will allow access to the ALF database identified by the g2config data group 199 located beneath the BASE_ALF_CONFIG entity 131.

Having described configuration of all services contained within the nodes of the server 13, configuration of the listener provided by the web server 12 is now described. As indicated above, a g2listener invocation parameter specifies an entity within the IDM which is used for listener configuration. Figure 28 shows an appropriate IDM data structure. In this case, the g2listener parameter is set such that:

g2listener=weblistener;

That is, the listener would locate a weblistener data group 210 located beneath a UAconfigs entity 211. The weblistener data group 210 indicates to where incoming requests should be directed (usually a CA service), and the name of a Java class providing authentication functions (in this case custom). The weblistener.custom data group 212 provides a mapping from request parameters to those needed for authentication.

It will be appreciated that means must be provided such that appropriate configuration data can be entered into the IDM database, to create the hierarchies described above. Specification of this configuration data is now described using a predetermined file format.

Figure 29 shows a file which can be used to configure the CL service 119 of the CA node 106. Text shown in square brackets “[]” denotes an entity names within the IDM hierarchy. Text shown in triangular brackets “< >” denotes a name of a data group which is positioned beneath the appropriate entity within the IDM hierarchy. Text shown in curly brackets “{ }” denotes a parameter type. “{GRF}” indicates that a parameter is a reference to another data group. A {GRF} parameter is followed by an appropriate entity name and data group name to which it refers. “{STR}” indicates a string parameter. Parameters specified by “{INT}” and “{BLN}” can also be used, although these are not shown in Figure 29. “{INT}” indicates an integer parameter, and {BLN} indicates a Boolean parameter.

Referring to Figure 29, it can be seen that the file specifies a CONFIG entity, having a SERVICE entity as a child, which in turn has a BASE_CL_SERVICE entity has a child. A CA_CL_SERVICE entity is a child of the BASE_CL_SERVICE entity, and has a LAB1_SrcApps entity as a child. The LAB1_SrcApps entity has a single g2config data group which comprises a single group reference entry which references a g2config data group located beneath a FirstByte entity within the hierarchical data structure.

The file of Figure 29 also specifies that the BASE_CL_SERVICE also has an EXT_APPS entity as a child entity, which in turn has a FirstByte entity as a child. The FirstByte entity has a single g2config data group which specifies two parameters of type string which provide authorization and protocol information relevant to the FirstByte application.

Having described configuration of the architecture shown in Figure 14, operation of services illustrated in Figure 14 is now described.

The UXM service 121 is responsible for receiving user requests, and generating one or more source application requests in response to the user request. The UXM service 121 is also responsible for creating composite pages.

It should be noted that the UXM uses an internal data format of UXMObjects to represent composite user interfaces and parts of such composite user interfaces. Using UXMObjects in this way allows the UXM service to operate in a manner independent of output format, and accordingly adds considerable flexibility to the system.

A UXMObject stores unstructured data, as well as attributes describing that data. A tree of UXMObjects is used to represent data for a composite user interface. It can be recalled that the structure of the composite application is described by UXM configuration data within the IDM database (see Figures 24 and 25). The tree of UXM objects will typically correspond to the tree structure for the UXM within the IDM, although it is important to note that the two trees are separate data structures which are stored separately.

A UXM object includes unstructured data for a particular entity within the UXM tree of the IDM database, which represents part of a composite user interface. This data is typically stored using an array of bytes. Each UXMObject also has a mode parameter and a type parameter. The mode parameter is used to indicate where within a composite document a UXMObject should be placed relative to other UXMObjects. The mode parameter can therefore take values such as before, after, replace, insert and toplevel. The type parameter is used to differentiate UXMObjects created by the UXM service

from UXMObjects created from source application data. UXMObjects can therefore have types of New and Existing. Each object additionally includes a set of attributes which are associated with the unstructured data, and which either describe the data, or are used for data conversion. Each UXM object has a identifier (which can conveniently correspond to an ID within the UXM tree within the IDM, see Figure 24), and references to parent and/or child objects if appropriate. These references to parent and/or child objects allow the UXMObjects to be used to create a UXM object tree. It should be noted that UXMObjects may also contain nested UXMObjects.

During operation, the UXM maintains a UXM tree (separate from the UXM object tree) which is based upon the configuration data stored in the IDM. This tree includes the actions, predicates and parameters which are present within the relevant entities and data groups of the IDM. Additionally, each entity within the UXM tree includes a state for each user at that node, which corresponds to a user's context.

The context for a user at a particular entity within the UXM tree will hold a state appropriate for that user for all requests which that user may make and all response which that user may receive. This is stored as a set of parameters (typically as NVPs), and by reference to a single object within the UXMObject tree.

In summary, it can be seen that the UXM essentially uses three data structures, the IDM data structure, the UXM tree and the UXMObject tree.

Information indicating how the UXM should operate in response to various requests and responses is configured within the IDM tree which at runtime is used to create the UXM tree. In addition to this information, the UXM tree also stored information relating to state of a particular user at a particular node is stored at a user's context at a particular entity within the UXM tree. Data pertaining to particular pages of the composite user interface is stored within a UXM object, and UXM objects collectively form a UXM object tree, which is separate from both the UXM tree and the IDM data structure. It can be deduced that for every entry in the UXM object tree, there must exist an entity in the UXM tree having the same identifier. The converse is not necessarily true.

Figure 30 shows a schematic illustration of part of a UXM tree 220 and a UXMObject tree 221. It can be seen that each entity in the UXM tree 220 includes context information for User 1. The context data associated with each entity in the UXM tree 220 includes reference (denoted by a broken line) to an appropriate instance of a UXMObject in the UXMObject tree 221. It can be seen that the UXM tree 220 and the UXM Object tree 221 have the same structure. It will be appreciated that in many embodiments of the present invention respective context data will be stored for a plurality of users, and each of the plurality of users will have a respective UXM object tree.

When describing the IDM data structure, it was explained that actions were executed if predicates were satisfied. Predicates which may appear in predicate data groups within the IDM are now described.

Some predicates are based on a user's log on credentials. Users may be allocated to one of a plurality of roles, and a RoleAuthorisedGuardCondition predicate entry specifying a particular role may be included in a predicate data group. This predicate will be satisfied only if a user's role matches the specified role. Similarly a RoleNotAuthorisedGuardCondition is satisfied only if a user's role does not match that specified by the predicate.

Some predicates are based on a NVP within a user's context at a given entity within the UXM tree. A CompareNVPValueGuardCondition predicate compares a specified NVP at a specified entity within the UXM tree with a specified NVP at a different entity within the UXM tree. The predicate has a parameter indicating whether it should evaluate to TRUE for equality or inequality.

A RegexNVPMatchGuardCondition predicate takes a regular expression as a parameter, and evaluates to true if a specified NVP at a specified entity within the UXM tree matches the regular expression.

An NVPEqualsGuardCondition predicate checks a value of a specified NVP in a user's context at the current entity against a value specified as a parameter, and evaluates to

TRUE in the case of equality. An `NVPNotEqualsGuardCondition` predicate performs the opposite function to an `NVPEqualsGuardCondition` predicate – that is it evaluates to TRUE in the case of inequality.

An `NVPExists` predicate checks if a specified NVP exists within the user's context at the current entity within the UXM tree. If the NVP does exist, the predicate evaluates to TRUE. An `NVPNotFound` fulfils the opposite function to an `NVPExists` predicate, that is, it evaluates to TRUE if a specified NVP does not exist at the current node.

Various other predicates can also be specified. A `RemoteServiceCredentialFoundGuardCondition` predicate determines whether a user's context at a specified entity within the UXM tree includes log on credentials for a specified external system (e.g. a source application). If the service credential exists, the predicate evaluates to TRUE. A `RemoteServiceCredentialNotFoundGuardCondition` predicate performs the opposite function, that is it evaluates to TRUE if a user's context at a specified node does not include a service credential for the specified external system.

`ExternalIDs` are used by the UXM to identify data elements within external applications, such as source applications. Some actions provided by the UXM are concerned with creating and using `ExternalIDs`. The UXM maintains mappings between `ExternalIDs`, and internal IDs used within the IDM.

An `ExternalIDFoundGuardCondition` predicate evaluates to TRUE if a user's context at a specified node includes a specified `ExternalID`. An `ExternalIDNotFoundGuardCondition` predicate evaluates to TRUE if the specified `ExternalID` is not found. A `NodeGuardCondition` takes a plurality of entity ID's as a parameter, and evaluates to TRUE if the ID of a specified entity matches one of the plurality of specified IDs.

A `DataReady` predicate checks if data (usually an `UXMObject`) is present at a specified entity within the UXM tree. This predicate is used to ensure that data necessary for aggregation (see below) is present at a given entity within the UXM tree.

A FalseGuardCondition predicate, and a TrueGuardCondition predicate may also be specified. A FalseGuardCondition always evaluates to FALSE effectively ensuring that the associated actions are never executed. A TrueGuardCondition always evaluates to TRUE ensuring that the associated actions are always executed.

A first task performed by the UXM service 121 is to generate requests to source applications following receipt of a request for a user. Operation of the UXM in performing this task is now described.

When a UserRequest message is received by the UXM service 121 (usually from the DT service 120), a UXM_MODIFIED parameter within the message is checked. If this is set to FALSE, the message is directed to a source application without processing by the UXM. If this is set to TRUE processing is carried out as follows and a UserRequest event is generated.

The UserRequest event causes the UXM to locate within the IDM the entity specified by the UXM_NODE_ID parameter. This entity then becomes a user's active aggregation point. In this case that node is the TreeNode 178. Any predicates specified in a predicate data group associated with that node are then checked. Assuming that all predicates are satisfied, the actions referenced by the entries of the integration data group (as defined by their respective data groups) are then carried out. The UXM traverses down through the UXM tree evaluating predicates and executing actions associated with each entity, each entity encountered typically providing a list of actions, one or more of which may reference an appropriate source application. Details of actions which may appear in an integration data group are set out below.

Actions are in general associated with a single event type. However, it is possible to override this associated by using an EVENTMASK parameter in an action's data group. This parameter is an N-bit binary value and is set such that each bit represents a different event, and that the action is associated with all events for which the respective bit is set to '1' – for example if four actions are used to trigger actions, the EVENTMASK parameter should be a four-bit binary value.

Actions typically associated with a UserRequest event are now described.

A UserRequestAction is by default associated with a UserRequest event, but may be associated with other events by using the EVENTMASK parameter in the manner described above. This action generates a new source application request which is targeted to the APPLICATION_CLASS of the source application. A request will be made to each source application which is required to produce the composite page.

The entries in an UserRequestAction data group specify information relevant to creating a request(s) for sending to particular source application(s), however it should be noted that details relating to connection to that application are not included, as these are specified by the CL service 123 of AA node 107. Each source application request will include only information specified by the relevant action data group, and will not necessarily include parameters present in the IncomingUserRequest message.

A source application request will include an APPLICATION_CLASS parameter which specifies an application class name for the source application, and an APPLICATION_PATH parameter, which specifies a path for locating the application. Additionally, a source application request may include a request method indicating how data should be requested. This parameter may take a value such as HTTP_GET in the case of a HTTP request. The message additionally includes a string comprising zero or more parameters which are picked up from the current context. If this parameter is the empty string, then no values are added to the source request.

A DeflectUXMRequest action is again associated with a UserRequest event by default, although this association can be overridden using an EVENTMASK parameter as described above. A DeflectUXM action generates a source application request of the type described above, but here the parameters for that request are obtained from the IncomingUserRequest message, not from the current context. This action is usually triggered to re-use an original user request message, and the UXM merely decomposes the message into parts relating to different pages of the source application.

A UXMRequestLink action is associated with a UserRequest event, and cannot be overridden by using the EVENTMASK parameter in the manner described above. This action triggers a UserRequest event on an entity within the UXM tree identified by a TARGETNODE parameter. The receiving entity will then respond to the UserRequest event as specified by its UXM entries. This action effectively adds an aggregation point to the UXM Tree.

A JumpToTreeNode action is again associated with a UserRequest event, but can be overridden using the EVENTMASK parameter. This action jumps to a different branch of the UXM tree structure, and causes a UserRequest event on an appropriate entity. This differs from the UXMRequestLink action described above in that the current aggregation point is removed, and replaced with the specified target aggregation point.

A number of actions are concerned with manipulating values (typically name, value pairs (NVPs)) within the current UXM context. A CopyNVPAction action is associated with a UserRequest event, but can be overridden using the EVENTMASK parameter as described above. A CopyNVPAction action copies sets an NVP at a specified target node to a value obtained from a specified source entity. The action includes identifiers of both the source and target entities, and details of the source NVP and target NVP within those entities. If a target NVP is not specified, as a default, it is assumed that the target NVP is the same as the source NVP. Optionally, a prefix and/or suffix may be specified which is prepended or appended to the source NVP when it is copied to the target NVP.

A ConcatenateNVPAction action is again associated with a UserRequest event but can be overridden by the EVENTMASK parameter. This action concatenates a plurality of NVPs at a source entity and copies these to a specified NVP at a target entity. The action takes as parameters a source entity identifier (which defaults to the current entity if not specified), a target entity identifier, a comma separated list of names from NVPs which are to be concatenated, and the name of an NVP to be created at the target entity. The value assigned to the created NVP will be the concatenation of the values from the NVPs in the comma separated name list.

An AddSequenceListener action is associated with a UserRequest event but can be overridden using the EVENTMASK parameter if necessary. This action sets the context of the current entity as a sequence listener of an entity specified by a TARGETNODEID parameter.

A StoreRemoteServiceCredential action is associated with a UserRequest event but can be overridden using the EVENTMASK parameter if necessary. This action allows information about a remote service (e.g. a source application) to be stored as a parameter in the context of the current entity. The information is stored as a NVP, and is typically security information related to a source application. The action includes a REMOTE_SERVICE parameter which is an identifier for a remote service (e.g. a source application), and a CRED_TYPE parameter which is used to indicate the type of credential being stored. The CRED_TYPE parameter can be set to USRPWD or IDONLY. If CRED_TYPE is set to USRPWD, the NVP stores both a username and a password. If it is set to IDONLY, only a username is stored. The name of the NVP is of the form “credential.<remote_service>.username|password”, where <remote_service> is an identifier obtained from the REMOTE_SERVICE parameter.

Having generated one or more source application requests in response to UserRequestAction actions and/or DeflectUXMRequest actions (in addition to carrying out any other actions which may be necessary), the created source application requests are used to create OutgoingUserRequest messages which are then usually targeted to a DT service 124 of an AA node.

A second function of the UXM is to compose responses received from source applications to form the composite application. The UXM service of the CA node generally receives source application responses from the DT service 124 of the AA node 107. An outline of the actions carried out by the DT service 124 of the AA node 107 is now presented, to aid understanding of the operation of the UXM.

The DT service 124 of the AA node 107 recognises pages returned by source applications using predefined rules. The recognition process is described in further detail below, but it should be noted that the DT attempts to recognise normal pages first,

then error pages. Assuming that the page is recognised a UXM object tree representing that page is generated.

Referring to Figure 31, there is illustrated a simple example conversion which may be carried out by the DT service 124 to generate an appropriate UXM object tree. A HTML document 222 comprises a table 223, which in turn comprises a form 224 and an image 225. The DT service generates a UXM object tree 226 from this HTML document. The UXM object tree 226 comprises four UXM objects, one for each element of the HTML document 222. It can be seen that these UXM objects are hierarchically arranged such that a ...Body object 227 is at the root of the tree, a ...Body.Table1 object 228 is a child of the ...Body object, and a ...Body.Table1.Form object 229 and a ...Body.Table1.Image1 object 230 are children of the ...Body.Table1 object 228. Thus, the hierarchy of the UXM object tree 226 mirrors that of the HTML document 222. It should be noted that each entry in the UXM object tree 226 is prefixed by “...” to indicate that in practical embodiments of the invention, the names of entries are prefixed by an appropriate source application name. Operation of the DT service 124 is described in further detail below.

Having created the plurality of UXM objects, these are sent to the UXM service 121 of the CA node 106.

A message is received by the UXM from the CA service of the AA node 107. If the message comprises one or more UXM objects, a UserResponse event is generated by the UXM. UXM objects may be nested within a received message, and in such a case the UXM unnests the received UXM objects to create individual UXM objects as appropriate. The created individual UXM objects are then copied to the user's context at the appropriate entity of the UXM tree structure and form the UXMObject tree. The appropriate entity can be identified using the ID parameter of the received UXMObject which, as described above, will correspond to the ID of an entity within the UXM tree.

The UserResponse event causes the UXM to traverse the UXM tree beginning at the entity specified by the ID parameter of the received UXMObject. For each entity within the tree which is processed in this way, the UXM determines if all predicates (specified

in appropriate predicate data groups) are satisfied for each entity. If the predicates are satisfied for a given entity, the actions associated with that entity (as specified in an actions data group) are executed. Details of actions which are generally associated with UserResponse events are now described.

A `RegexTransformNVPValue` action is by default associated with a UserResponse event, although this association can be overridden using the `EVENTMASK` parameter as described above. This action sets a NVP within the user's context, for a target entity within the UXM tree. The action specifies a `SOURCENODEID` parameter which defaults to the current entity if not specified, and a mandatory `SOURCENAME` parameter which specifies an NVP within the source entity. The action may optionally specify `TARGETNODEID` parameter identifying a target entity and a `TARGETNAME` parameter identifying a NVP within the target entity. If these are not specified, the `TARGETNODEID` is set to be equal to the `SOURCENODEID` and/or the `TARGETNAME` is set to be equal to the `SOURCENAME`. The action also includes a `REGEX` parameter which specifies a regular expression which is applied to the source NVP to generate the target NVP.

Regular expressions will be well known to those skilled in the art as a way of manipulating strings within computer programs. In the case of some embodiments of the present invention regular expressions of the type used in the Perl5 programming language are preferably used, although it will be readily apparent to those skilled in the art that regular expressions of different formats can also be used in embodiments of the present invention.

A `CopyUXMObjetValueToTargetContext` is again associated with a UserResponse event, but can again be overridden using the `EVENTMASK` parameter as described above. This action sets an NVP in a user's context at target entity within the UXM tree. The value to which the NVP is set is taken from a UXM object at a source entity. The action takes `SOURCENODEID` and `TARGETNODEID` parameters which respectively specify source and target entities within the UXM tree. If these parameters are not set, by default the current entity is used. The action has a mandatory `TARGETNAME` parameter which specifies an NVP within the target node which is to be set.

A SetUXMObjectProperties action is again associated with a UserResponse event, but can be overridden as described above. The action specifies a plurality of NVPs which are set within the UXM object at the current entity of the UXM tree.

A SetTargetUXMNode action is associated with a UserResponse event, but can be overridden using the EVENTMASK parameter. This action sets attributes of a UXMObject at an entity specified by an ACTIONTARGET parameter, which defaults to the current entity if no value is specified. The UXM object specified by the ACTIONTARGET parameter is updated so that it points to a UXMTreeNode relating to a different entity within the UXM object tree specified by a mandatory TARGETNODE parameter.

A CopyUXMObjectToTargetContext action is again associated with a UserResponse event but can be overridden. This actions sets the UXMObject at an entity specified by a SOURCENODEID parameter to point to the UXMObjects specified by an entity given in a TARGETNODEID parameter. The SOURCENODEID parameter is optional, and defaults to the current entity if not specified.

A RelocateUXMObject action is associated with a UserResponse event and cannot be overridden using the EVENTMASK parameter. This action sets the mode parameter within a UXM object associated with an entity specified by a TARGENODEID to "Relocate", and results in the object being relocated (i.e. moved rather than copied) to a new parent.

A DeleteUXMObjectFromTargetContext action is again associated with a UserResponse event but can be overridden. This action deletes the UXM object associated with the entity specified by a mandatory TARGETNODEID parameter.

A ConditionalCopyOfUXMChildObjectToTargetContext action is associated with a UserResponse event but can be overridden. This action compares a value of an NVP specified by a SOURCENVPPNAME within the context of the current entity with the same NVP in an entity referenced by a CHILDNODETOCOMPARE parameter. If the

comparison is such that a predetermined condition is satisfied, a UXM object identified by a CHILDNODETOCOPY parameter is copied to the context of an entity identified by a TARGETNODEID parameter. It should be noted that the parameters CHILDNODETOCOMPARE and CHILDNODETOCOPY are specified relative to the current entity, and not as absolute path names.

A CopyUXMObjectValue action is again associated with a UserResponse event and cannot be overridden using the EVENTMASK parameter. This action copies values from a UXM object present at an entity identified by a SOURCENODEID parameter to a UXM object present at an entity identified by a TARGETNODEID parameter.

A CopyUXMObjectValueFromTargetContext is again associated with a UserResponse event but can be overridden. This action will copy a value of an NVP identified by a SOURCENVPPNAME parameter in the context of an entity identified by a SOURCENODEID to a UXM object at an entity identified by a TARGETNODEID parameter.

A RegexTransformUXMObjectValue is by default associated with a UserResponse event but this can be overridden using the EVENTMASK parameter. This action retrieves a UXM object from an entity identified by a SOURCENODEID parameter, applies a regular expression specified by a REGEX parameter and writes the resulting value to an entity identified by a TARGETNODEID parameter.

A UXMObjectNotification action is associated with a UserResponse event and cannot be overridden. It takes no parameters. It copies a UXMObject from a current context to the child of an active aggregation point having the same name. For example, if the action is associated with an entity having a path S1.A1.A2 and the active context is S1.C1, the id of the current entity (i.e. A2) is used to locate an entity within the active context which should refer to the UXM object. In this case, that is S1.C1.A2.

Those actions which relate to manipulation of ExternalIDs and which are by default associated within a UserResponse action are now described.

A LoadExternalID action can have its association overridden using the EVENTMASK parameter described above. This action retrieves an ExternalID which is specified by an EXTERNAENTITYCLASS parameter and an APPLICATION_CLASS parameter from the current context, and writes this value to a UXM object associated with the current context.

A StoreExternalID action can again have its association overridden by use of the EVENTMASK parameter. This action saves the value of the UXM object associated with the current entity as an ExternalId having an application class identified by an EXTERNAENTITYCLASS parameter.

A ClearExternalIDs actions can be overridden, and clears all external ID's from the context of an entity identified by a SOURCENODEID parameter.

Having traversed the tree, and executed all actions for which predicates are satisfied, execution returns to the user's active aggregation point (i.e. the entity within the UXM tree which received a UserRequest event to trigger the source application requests which in turn generated the response from the AA). Each node of the UXM tree beneath the node which received the UserRequest event is interrogated to ascertain whether or not it has received its UXM object. When all entities beneath the entity which received the request have received their objects, a UXMAggregate event is created to cause aggregation.

It should be noted that in preferred embodiments of the present invention, each node in the UXM tree can specify that a UXM object is mandatory or optional within its control data group using a Boolean isMandatory variable, which defaults to FALSE where not specified. In such embodiments, a UXMAggregate event is created as soon as all mandatory UXMObjects are present. However, in alternative embodiments of the invention it may be desired to await all mandatory UXM objects, then apply a predetermined time out, after which aggregation is carried out using those objects which have been received.

A UXMAggregate event causes the UXM to traverse all entities which are children of the current entity, and execute all actions specified in the integration data group which are triggered by a UXMAggregate event. Some actions which are associated with UXMAggregate events are now described.

An AddUXMObjectAction action is by default associated with a UXMAggregate event, but this can be overridden using the EVENTMASK parameter described above. This action creates a new UXMObject and adds this to an entity identified by a TARGETNODEID parameter. This action therefore allows UXMObjects to be created at any point in the UXM tree. An OBJECTID parameter specifies a name for the new UXMObject, and OBJECTTYPE and MODE parameters specify the created object's type and mode as described above. A RELATIVETO parameter allows the TARGETNODEID to be specified relative to an entity specified by the RELATIVETO parameter, although if RELATIVETO is not specified, it defaults to the top level.

As described above, the AddUXMObjectAction action creates a new UXM object at a target entity. An INSIDETARGET parameter allows configuration of what should happen if a UXM object already exists at the target entity. If the parameter is not specified, and the UXM object already at the target entity is of type "container" the current object is provided as a child of the existing UXM object at the target entity. If the INSIDETARGET parameter is set to TRUE, the new UXMObject is set as a child of the existing UXM object. If the INSIDETARGET parameter is set to FALSE, the existing UXM object becomes a child of the new UXM object at the target entity.

A CreateUXMObject action is by default associated with a UXMAggregate event, but this can be overridden using the EVENTMASK parameter as described above. This action creates a new UXMObject and adds it to the user's context at the targeted entity. It should be noted that everything which can be achieved using a CreateUXMObject action can be created by an AddUXMObjectAction which has appropriate parameters. That is the CreateUXMObject action can be thought of as an AddUXMObjectAction with some hard-coded parameter values.

An AggregateChildObjects action is by default associated with a UXMAggregate event. This cannot be overridden using the EVENTMASK parameter described above. This action is essential to creating composite user interfaces. It aggregates all UXMObjects at contexts of entities which are children of a target entity, and creates a new UXMObject at the targeted context. Any UXMObjects which have a mode set to delete are ignored for the purposes of this action, all others are inserted as nexted UXMObjects. A parameter for the action specifies the target entity.

An AggregateNamedContexts action is by default associated with a UXMAggregate event, but this can be overridden using the EVENTMASK parameter described above. This action allows zero or more entities to be specified, and the UXMObjects associated with each of those entities are aggregated to form a new UXMObject at an entity specified by a TARGETNODEID parameter.

A DeleteUXMObjectAction action is by default associated with a UXMAggregate event, but this can be overridden using the EVENTMASK parameter described above. This action will set any UXMObjects present at an entity defined by a TARGETNODEID parameter to have a Mode of Delete, meaning that it will not be included in aggregation actions (see above). It should be noted that this action does not delete objects, but simply sets their mode parameter to have a value delete.

An AddHTMLElementAction action is by default associated with a UXMAggregate event, but this can be overridden using the EVENTMASK parameter described above. This action is used to inset a HTML element into a UXMObject in the context of an entity identified by a TARGETNODEID parameter. In order for this action to work, the UXMObject must have a type parameter as specified by an ENUMERATED_HTML_ELEMENT_TYPE parameter within the AddHTMLElementAction action. Possible values for the ENUMERATED_HTML_ELEMENT_TYPE parameter are html_radiobuttongroup, html_checkboxgroup and html_dropdownlist. An AddHTMLElementAction also takes parameters which specify a name, a value and text for the HTML element. For each type of HTML element further parameters relevant to that element may be specified by the AddHTMLElementAction.

If a UXMObject at a target context has a type of `html_radiobuttongroup`, an `AddHTMLRadioButton` action is by default associated with a UXMAggregate event, but this can be overridden using the `EVENTMASK` parameter described above. This action adds a radio button to the radio button group represented by the UXMObject.

Having executed all actions appropriate to a UXMAggregate event, the entity within the UXM tree which received the user's request has a UXMObject which represents the requested composite page, which can be passed to the DT service 120 of the CA node 106.

It should be noted that UXMObjects are passed between services by creating an XML representation which is then used as a payload of a message to be transferred between services. This is achieved using a `toXML` method provided by a UXMObject.

It should be noted that the UXM comprises further functionality in addition to that which is described above. Actions triggered by various events have been described above. It should be noted that some actions may by default be associated with no particular event, and in such cases the `EVENTMASK` parameter must be used to determine association.

The UXM service 121 allows JavaScript scripts to be associated with given entities within the tree, and these scripts can be executed by inclusion in an action data group, in the manner described above. The JavaScript is interpreted by the UXM at run time, and executed to perform the necessary functions. It will be appreciated that scripts may need to communicate with the UXM to obtain relevant data. This is achieved by providing a `UXMRUNTIME` object which can be accessed by scripts.

Methods exposed by a Java object can be invoked directly from a script, assuming only that the script has the Java object's identifier. The methods exposed by a `UXMRUNTIME` object allow scripts to be effectively created.

A UXMRUNTIME object allows various operations to be carried out on UXM objects specified by a path within the UXM tree. This path is a string comprising the name of an entity prepended by a '.' prepended by the name of its parent entity, which in turn is prepended by a '.' and its parent entity name. For example the string "MyTRN.page1.FormElement" denotes an entity "FormElement" which has as its parent an entity "page1" which is a child of a TreeRootNode entity "MyTRN". It should be noted that fully specified paths must always be used within UXM scripts. This is because a UXM script is not considered to be executed from a particular entity within the UXM tree, and accordingly a relative path will have no well defined meaning.

Methods provided by the UXMRUNTIME object are now described.

A prepareUXMObject method takes a path of the form described above as a parameter and allows the UXMObject specified by the path to be accessed by the script. The thread in which the script is executing is blocked until the UXMObject is returned, and the UXMObject's availability can therefore be assumed. This method will effectively trigger a UserRequest event on the node specified by the path parameter, and it is therefore important to ensure that nodes which are to be used in this way by scripts are provided with actions triggered by UserRequest events. An entity targeted by this method is typically a source page, and this method effectively creates a request to fetch the source page, which is subsequently provided to the script. If the UXMObject cannot be provided to the script for any reason, an error field within the UXMRUNTIME object will contain details of any exception which was thrown. The error field of the UXMRUNTIME object should therefore be checked (see below) before attempting to use the requested UXMObject. A prepareUXMObjects method functions analogously to prepareUXMObject, but takes an array of paths as its parameter, and returns all requested UXMObjects.

A getUXMObject method retrieves a UXMObject from an entity specified by a path parameter. A setUXMObject method sets a UXM object at a specified entity to a UXMObject provided by a parameter of the method.

The UXMRUNTIME object provides various methods for accessing and manipulating NVPs at a given entity within the UXM tree. A setNVP method allows a specified NVP to be set to a specified value at a specified entity within the UXM tree. A getNVP method allows the value of a specified NVP to be obtained from a specified entity. A clearNVPs method allows all NVPs to be cleared from a specified entity. A getRequestParameter returns the value of a request parameter at a specified entity in the UXM tree.

It was mentioned above that the UXMRUNTIME object includes a field in which details of any errors are stored. A hasException method returns a value indicating whether or not an exception occurred during processing, and a getException method returns details of any exception that occurred.

Error details may also be stored within a user's context at an entity of the UXM tree. Various methods are provided to access such error details. A getErrorDetail method gets details of such an error message from an entity specified by a path parameter. A requestHadError method returns a value indicating whether or not an error occurred. A getErrorLocation returns the name of a service (e.g. UXM, DT, CL) at which the error occurred.

A getErrorMessage method returns details of the error message generated, and getMajorErrorCode and getMinorErrorID methods return details of error codes. All these methods take a path specifying an entity within the UXM tree as a parameter.

In general terms, scripts can be used to replace the mechanisms described above for carrying out aggregation. When a script is used an integration data group will typically contain a single entry referencing an action data group for the script within an action library which contains a script name and details of the script.

If a script is used for aggregation, it will be appreciated that it is important to ensure that a mechanism is provided such that the script knows when all necessary source pages have been received by the UXM as UXMObjects. This is provided by a ScriptResponseNotification action, which is provided in an action library, and included

in the integration data group of each source page, and which is triggered after processing of a UserRequest event.

The runtime memory space of the script will include details of all requested source pages. When a UXMObject corresponding to a source page is received by the UXM the UXMRUNTIME object is updated by the ScriptResponseNotification. The UXMRUNTIME object in response to the ScriptResponseNotification then updates the runtime memory space of the script.

It should be noted that when scripts are used, two parameters which may be specified in the control data group associated with an entity are of some importance. An EVENTORDER parameter controls the order in which child entities are processed. A HANDLEERROR parameter needs to be set to TRUE, so as to allow the script to process any errors that may occur during execution, instead of the global error handling generally provided by the composer.

The scripting functions described above are provided by the Bean Scripting Framework (BSF) which is provided by the Java programming language. This framework allows a number of different scripting languages to be used, although as described above, JavaScript is used in preferred embodiments of the present invention. It should be noted that on first execution the JavaScript script will be compiled to generate Java Byte Code which can be interpreted by a Java Virtual Machine. It will therefore be appreciated that a script will take longer to execute on its first execution.

Operation of the DT service is now described. The DT service includes a plurality of transformation engines which may be applied to data to cause data transformations. The transformation engines are stateless, and a single instance of each engine exists within a particular DT service. These transformation engines are now described in further detail.

An APICallExternalizer translates data from internal representations to external representations using ExternalIDs described above and such data can then be used by the CL service to generate API calls. A DOMtoString transformation engine converts a

DOM object representing an HTML file (typically received from a source application) into a string.

An IDExternaliser transformation engine and an IDInternalizer transformation engine respectively convert internal IDs to external IDs and external IDs to internal IDs.

The DT service includes some transformation engines to process and correct HTML provided to it. A HTML TagBalancer tag balances incoming HTML to produce well formed HTML as output. A HTMLTidy transformation engine makes use of Jtidy to correct malformed HTML, and ensures that the output is XHTML compliant.

A JavaSpecializedEngineExecutor transformation engine executes specialized transformations that are hardcoded into a Java class. The Java class specified for this engine must implement a predefined JavaEngine interface specifying a single execute() method which is used to transform data, so that it can be known that the Java class will include this method which is required for transformation.

A Marshaller transformation engine makes use of the Castor framework to convert Java objects into XML, as described at <http://castor.exolab.org>. This engine can be used for generic transformations.

A PageResolver transformation engine recognizes response pages and applies appropriate extraction configuration. It first tries to recognize normal pages, then error pages and finally applies a default extraction if no match is found.

A RegEx transformation engine takes a string input, applies a regular expression to it and returns a string as output.

A RegexAggregator transformation engine is used to insert new user interface elements into a user interface. The configuration of this engine is a single regular expression string. Some regular expressions have special meanings, and will be substituted before the regular expression is applied:

__VALUE__ is replaced by the value of the OBJECT_ID parameter in the engine definition data-group.

__NAME__ is replaced by the value of the NAME parameter in the engine definition data-group.

A RegexExtractor transformation engine is used to extract user interface elements in UXM format using regular expressions. It also marks up the input string where the string fragment was extracted.

A UIAggregator transformation engine is used to aggregate the UXM user interface data into a form suitable to return to the user. It is typically only used in the CA.DT. A UIExtractor transformation engine converts the incoming response page into UXM Objects which can be forwarded to the UXM service. This is used in the AA.DT.

An Unmarshaller transformation engine makes use of the Castor framework to convert XML into Java. A UxmObjectDecomposer transformation engine is used in the process of aggregating UXM user interface data. It decomposes complex UXM objects into simpler ones so that the aggregation can take place in later stages. A UxmObjectToString extracts the data from a UXMObject and returns it as a string.

An XMLParser transformation engine takes a string as input and parses it into a DOM tree. The resulting XML tree is the output of the engine. An XPathExtractor transformation engine is used to extract user interface elements in UXM format using XPath. It also marks up the input string where the string fragment was extracted. It is used primarily in the UXM. An XSLTransformer transformation engine takes an XML document as input, applies a XSL style sheet to it and returns an XML document as output.

The transformation engines outlined above provide a variety of basic functions which are required by a DT service. The transformation engines are used by transformation actors which act on the data to transform it. There are essentially two types of transformation actors: atomic actors make use of transformation engines directly, while group actors use other transformation actors.

An atomic actor holds configuration data indicating how a transformation engine should be used to carry out the necessary transformation. For example, an atomic actor using an XSL transformation engine may be configured using an XSL style sheet. The output will then be a DOM object (see above). When called, the actor will pass the input data and the configuration data to the transformation engine to cause the transformation to be effected.

An atomic actor has a type parameter which is set to "base" which indicates that the transformation actor is an atomic actor. It has a textual description of the transformations which it carries out. It has a EngineID parameter which identifies one of the transformation engines set out above which is to be used by the actor. A configuration parameter is represented by a string. This will typically be a regular expression or alternatively a file name of an XSL style sheet. Some actors will have multiple outputs, while other actors will always have a single output value. This is represented by a MultipleOutputs parameter which is set to TRUE if multiple outputs are allowed, or FALSE if MultipleOutputs are not allowed. If an actor for which MultipleOutputs is set to FALSE produces multiple output data at run time, the data is merged using a Java class specified by a DataMerger parameter. A Timeout parameter provides a timeout, preferably in milliseconds.

A group actor combines one or more other transformation actors. The transformation actors being grouped may be atomic actors or group actors, and can be executed either in parallel or sequentially, or as alternatives.

A group actor has a type parameter which is set to "group", a description parameter specifying a textual description, and a parameter TransformationIDs which is a list of actors contained within the group actor. A GroupType parameter can take values of "sequence", "parallel" and "switch". If the GroupType parameter is set to "sequence" all the contained transformation actors are executed sequentially, in the order in which they are listed in a Transformation IDs parameter. If the GroupType parameter is set to "parallel", all the contained transformation actors are executed in parallel. If the GroupType parameter is set to switch, a single actor will be chosen at run time to act

upon the passed input data. A LateInitialization parameter takes a Boolean value indicating whether or not late initialisation should be used. That is, this parameter indicates whether the transformation engine should be initialised at start up (if LateInitialization is set to FALSE), or when the transformation engine is used for the first time (if LateInitialization is set to TRUE). A group actor also has MultipleOutputs, DataMerger and Timeout parameters, as described above with reference to atomic actors.

In preferred embodiments of the present invention some atomic actors and some group actors are provided by default. Details of these transformation actors are now described.

Atomic actors provided by default are as follows. A NOP actor performs no operation, which can be necessary in the DT in cases where no transformation is required. Three further atomic actors simply make transformation engines described above available as actors. A HTMLTagBalancer actor makes the HTMLTagBalancer transformation engine available as an actor, a UXMObjectDecomposer actor is used by the DT service of the CA and makes the UXMObjectDecomposer transformation engine available as an actor. A UXMObjectToString actor makes the UXMObjectToString transformation engine available as an actor, and an XMLParser actor makes the XMLParser transformation engine available as an actor.

A PageResolver actor is used by the DT service of the AA to recognise pages returned by source applications. It picks up details of the pages to be recognised from the appropriate parts of the IDM, as described above. A UIExtractor actor is used by the DT service of the AA to convert pages into UXMObjects, again appropriate configuration information is provided by the IDM.

A UIAggregator actor is used by the DT service of the CA to convert UXM objects into a form suitable for returning to a user as part of a composite application. A URLRewrite actor is used to update any necessary URLs, forms and/or frames in a page which is to be returned to a user.

A plurality of group actors are also provided by default. A sequence.IncomingUserResponse actor is a sequence of transformation actors which can be applied to a typical incoming user response message, to convert the page response into UXM form. It applies the PageResolver actor and the UIExtractor actor described above.

A sequence.OutgoingUserResponse actor is sequence of transformation actors to apply to a typical outgoing user response message to convert the UXM Object tree into the resulting page to return to the user. It applies an actor to convert a marshalled UXMObject tree into a XML tree and then aggregates the various components according to information in the XML tree. The result is converted to a string representing the page which is returned to the user.

A sequence.UIAggregation actor comprises a sequence of transformation actors to apply to aggregate the components of a page into the final page. It applies other transformationa actors to decomposes and aggregate and then uses the atomic actor UIAggregator described above.

A sequence.UIAggregation.DecomposeAndAggregate is a sequence of transformation actors which are applied to get the page elements ready for aggregation. It first decomposes the top most UXMObject using the UXMObjectDecomposer actor. It then recursively goes down through the tree repeating the same action.

A switchgroup.UIAggregationRecursion is a switch group of transformation actors. It goes down through the UXMObject hierarchy, choosing at each stage whether the UXMObject contains other UXMObjects or not.

A sequence.UIExtraction actor comprises a sequence of transformation actors which are applied to extract the components of a page into a UXMObject tree. It decomposes and aggregates, then uses the atomic actor UIAggregator

A switchgroup.UIExtractionRecursion actor is a switch group of transformation actors. It goes down through the UXMObject hierarchy, choosing at each stage whether the UXMObject contains other UXMObjects or not.

Operation of the DT service within the AA node for a particular source application is illustrated in Figure 32. Data is received from a source application is denoted by an arrow 231. This data is first passed to a transformation engine 232 which makes any necessary amendments to URL references included in the received source data. This can be achieved using, for example, the RegEx transformation engine described above. A second transformation engine 233 performs any generic URL amendments that may be necessary, and can again be the RegEx transformation engine described above. A HTML tag balancer transformation engine 234 (of the type described above) is then used to ensure that the produced HTML is well formed.

The output of the HTML tag balancer is input to a group actor 235 which is configured to recognise and process the received source page. The group actor 235 comprises a suitable PageResolver transformation engine 236 which is configured to recognise predefined pages, and a group actor 237 which comprises one or more of the transformation engines described above, which together act to extract user interface elements from recognised pages.

Having described configuration of the composer, and the key services contained within its nodes, its operation in providing composite applications, as outlined above with reference to Figure 15, is now described in further detail, referring again to Figures 14 and 15.

Processing typically begins when a user using the web browser 126 makes a HTTP request which is subsequently directed to the web server 12. The servlet 117 operating on the web server 12 determines whether or not the URL entered relates to a page provided by the composer. If it is determined that the URL does not involve the composer, the web server obtains and displays the requested web page in a conventional manner. If the URL is to be provided by the composer, the servlet 117 intercepts and

processes the HTTP request. This involves converting the HTTP request into an IncomingUserRequest JMS message.

It will be appreciated that some requests will require authentication. Authentication is handled using cookies. When a user first attempts to access a page requiring authentication the request will be intercepted by the CL service of the CA node, and this service will determine whether or not the request includes any required cookie. If no such cookie is present in the request, the user will be presented with a page into which suitable authentication data (e.g. username and password) is entered. This data is sent to the web server, and onwards to the ALF service 126 of the ALF node 108, where the input authentication data is checked against data stored in the ALF database. Assuming that this authentication process is successful, the user's original request is processed. A cookie is stored on the user's computer, and in the case of subsequent requests, this cookie is read by the web server, forwarded to the CL service of the CA node, and then forwarded to the ALF service 126 to enable authentication using the ALF database 110. The following description assumes that all necessary authentication has been successfully carried out.

Creation of the JMS message is carried out as specified by the appropriate configuration data within the IDM. Having created an appropriate IncomingUserRequest message, and determined to where it should be sent using the target parameter within the weblistener data group 210 (Figure 28) it is then necessary to locate an appropriate instance of the target service which can accept the created message.

The listener performs load balancing over JMS, so that a plurality of instances of a target service can be used concurrently, and requests can be effectively distributed between the plurality of composers. The JMS properties file for the listener (described above) includes a resolverprotocol.roundrobin.target parameter which specifies a load balancing algorithm which is to be used such as, for example the round robin algorithm. This parameter will also specify how many instances of the target (in this case the CL service of a CA) to expect.

When an instance of the target service starts up, it declares itself using an appropriate JMS message, and the listener will then know that messages can be sent to that instance of the target service. When a listener needs to send a message to an instance of the target service (e.g. to the CL service of a CA) a broadcast message is sent to all instances of the service registered with the listener. The listener will then await a response from all registered targets, any responses from non-registered targets are simply ignored. When all expected responses have been received, the listener chooses one instance of the target service in accordance with the specified algorithm, and the IncomingRequestMessage is sent to the selected instance.

Having selected a target to receive messages from a particular user, the listener ensures that subsequent messages from that user are directed to the same composer. This need not necessarily be the case, and is determined by an add_affinity parameter within the IDM. Similarly, the IDM configuration for the composer has an add_listener_affinity parameter which ensures that messages for a particular user are always directed to the same listener.

The transmitted IncomingUserRequest JMS message will be received by the CL service 199 of the CA node 106 via its ML service 118 which is responsible for JMS messaging. On receipt of the IncomingUserRequest message, the CL may perform authentication using log on parameters specified in the IncomingUserRequest message, by using the ALF service of the ALF node in the manner described above. Assuming that authentication is successful, the message is simply passed to the DT service 120 of the CA node. The passage of messages between services, is determined by data stored within the message objects, as described above.

The DT service 120 retrieves configuration data pertinent to the actions required in response to receipt of an IncomingUserRequest message. The data is retrieved by accessing the IDM data structure (Figure 22), and locating the DT.Applications data group 153. The data group 154 pertinent to the composite application specified in the IncomingUserRequest message is then located, and the action referred 158 referenced by the IncomingUserRequest message in the data group 154 is then identified and carried out. On receiving an IncomingUserRequest message, no transformation is

usually required, and the DT service 120 therefore performs no operation, as specified by the data group 158. The message is then passed to the UXM service 121 of the CA node 106.

The IncomingUserRequest message is typically generated from a HTTP request as described above. A suitable HTML fragment for such a request is shown in Figure 33. It can be seen that the third line specifies that the composite application's application class name is "CompApp1". The fourth line specifies that the request requires modification by the UXM, the fifth line specifies a node in the IDM tree structure relevant to the application, and the sixth line specifies a node in the IDM relevant to the page which generated the request. All this data is encapsulated within the IncomingUserRequest message.

The UXM will locate a "CompApp1" node within the UXM tree, and traverse down the tree structure, executing any actions for which the predicates are satisfied. These actions will typically involve creating one or more OutgoingUserRequest messages directed to appropriate source applications, which are forwarded to the source applications via the AA node 122, and these messages are therefore forwarded to the AA node.

In the present example, an AA node is present within the same process as the CA, and can therefore be located without problems. However, if no such node exists within the current process. The message(s) are forwarded to the ML service 118 of the CA node 106 and the ML service then attempts to locate an appropriate AA node within a different process using JMS, as described above with reference to Figure 12D. When a suitable node has been located, the message(s) are forwarded to the ML service of that node, and onwards to the CL service of that node. Using JMS in this way allows convenient inter-process communication, regardless of the distribution of processes between physical machines.

An OutgoingUserRequest message is received by the DT service 124 of the AA node 107. The DT service must transform the OutgoingUserRequest into a form understandable by the target source application. In many cases, there will be no action required by the DT at this stage as shown in the example configuration of Figure 23,

where it can be seen that the OutgoingUserRequest entry of the DT.App.firstbyte data group 161 references a no operation action 162. The OutgoingUserRequest message is therefore forwarded to the CL service 123 of the AA node 107.

The CL service 123 uses configuration data stored within the IDM to locate a data group representing the correct source application (e.g. the g2config data groups 145, 147 of Figure 21). This data group will specify any authentication information which must be included in a request to that source application. In order to obtain such authentication information the CL service 123 requests authentication data from the ALF service 126 of the ALFNODE 108, and this data is then fetched from the ALF database as indicated by data in the IDM.

As described above, the IDM configuration for a CL service also specifies parameters which are used to connect to the appropriate source application (specified in a connection data group). The CL service will use these parameters to determine the format of request which should be generated for onward transmission to the source application. For example, this may be a HTTP request directed to a particular port of a particular host. This request is then transmitted to the source application.

The source application will process received request messages in a conventional manner, and data output in response to the request is in turn received by the CL service 123 of the AA node 107.

The CL service 123 receives the data from the source application, and forwards this data to the DT service 124 of the AA node 107. At this stage the DT service 124 must act upon the received data to transform it into a form suitable for sending to the UXM service 121 of the CA node 106. This processing will be of the type illustrated in Figure 32 above.

The created UXMObjects are then forwarded to the UXM service 121 of the CA node 106. The UXM will record the received objects within the UXM Tree as described above, and when all necessary data has been received, aggregation will occur to create one or more UXMObjects representing the requested composite page. The created

UXMObjects are then forwarded to the DT service 120 of the CA node 106, in one or more messages containing XML versions of the created UXMObjects.

The DT service 120 of the CA node 106 must then receive the messages, recreate the UXMObjects, and use data from the IDM configuration to determine how to convert the received UXMObjects into a form suitable for output to the user. Suitable forms may include HTML, or plain text.

It should be noted that by using an internal form until the composite page is processed by the DT service 120 of the CA node 106, the system described above can easily be adapted to handle different output types, simply by adding appropriate configuration data to the IDM for the DT service 120. This allows an effective decoupling of the composer from the output format required. Similar decoupling can be provided when handling input data.

The created output data is then forwarded by the DT service 120 to the CL service 119. The CL service uses the IDM data structure to determine the protocol which should be used to return the composite page to the user, usually via the servlet 117 of the web server 12.

Referring back to Figure 14, operation of the administration terminal 112 which is connected to the service 13 by means of the connection 113 is now described. It is preferred that the administration terminal 112 is connected to the server by a network connection, and a conventional Transmission Control Protocol/Internet Protocol (TCP/IP) connection of the type used in Internet networks can suitably be used. The administration terminal may operate by running a web browser and accessing an adminconsole HTML document which is provided on a suitable port number of the server (e.g. port 8080).

When a user has logged onto the administration server, it is possible to create users, and roles. Users are allocated to roles, and the roles to which a user is allocated determine the actions which may be performed by that user. Roles specify commands which may be executed by their users and may also specify child nodes which inherit their properties.

When roles have been created users can be allocated to created roles. Users can be removed from roles as necessary, and deleted if no longer required. If a role is deleted, data for all users allocated to that role is automatically updated. The administration terminal can also be used to set up a user's service credentials, that is the security data needed by a user to access various source applications used within the composite applications.

Although the embodiment of the invention described herein is implemented using the Java programming language it will be appreciated that the invention is in no way restricted to a Java implementation. Indeed, the invention can be implemented using any appropriate high level computer programming language. However, it is preferable to use an object oriented programming language to obtain the benefits of reuse and encapsulation which are inherent in such computer programming techniques.

Although the present invention has been described hereinbefore with reference to particular embodiments, it will be apparent to a skilled person in the art that modifications lie within the spirit and scope of the present invention.